



An Alexander polynomial for MOY graphs

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Abstract

We introduce an Alexander polynomial for MOY graphs. For a framed trivalent MOY graph \mathbb{G} , we refine the construction and obtain a framed ambient isotopy invariant $\Delta_{(\mathbb{G},c)}(t)$. The invariant $\Delta_{(\mathbb{G},c)}(t)$ satisfies a series of relations, which we call MOY-type relations, and conversely these relations determine $\Delta_{(\mathbb{G},c)}(t)$. Using them we provide a graphical definition of the Alexander polynomial of a link. Finally, we discuss some properties and applications of our invariants.

Keywords Alexander polynomial · MOY graph · State sum · MOY-type relations

Mathematics Subject Classification Primary 57M27 · 57M25

1 Introduction

1.1 Main results

We study a class of spatial graphs, which we call MOY graphs. The terminology was first used by Wu [14] in his categorification of the MOY calculus introduced by Murakami-Ohtsuki-Yamada [11]. Roughly speaking, an MOY graph is an oriented spatial graph in the 3-sphere S^3 satisfying the following conditions:

- (i) *Transverse orientation* For every vertex v , there is a straight line L_v that separates the incoming and outgoing edges.

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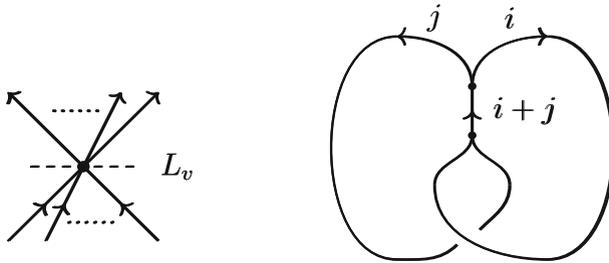


Fig. 1 The local picture of a vertex for an MOY graph (left) and an example of MOY graph (right)

- (ii) *Balanced coloring* Each edge is colored by a non-negative integer such that for every vertex v , the sum of the colors of the edges entering v equals that on the edges leaving v .

While we will give the precise definition in Sect. 2, we exhibit an example of MOY graph in Fig. 1. We remark that an oriented spatial graph with a transverse orientation is called *transverse spatial graph* in Harvey-O’Donnol [5, Definition 2.2], where an alternative definition using the formal description of simplicial complex was given.

We then introduce the Kauffman state sum $\langle D, c \rangle$ for a given graph diagram D of G with a coloring c , which generalizes the Kauffman state sum introduced in [6] for a knot. Based on the state sum we define an *Alexander polynomial*

$$\Delta_{(G,c)}(t) := \frac{\langle D, c \rangle}{(t^{-1/2} - t^{1/2})^{|V|-1}},$$

where $|V|$ is the number of vertices of G . The Alexander polynomial is a rational function of $t^{1/2}$ and $t^{-1/2}$, and is, like its classical counterpart for a link, well-defined up to t^k .

When we fix a framing on G and obtain a framed graph \mathbb{G} , it is possible to construct a *normalized Alexander polynomial* of \mathbb{G} that eliminates the aforementioned ambiguity. For the reader who is familiar with the definition of Jones polynomial of knots, this is the same idea as multiplying a writhe factor to Kauffman’s bracket polynomial. We will focus on the study of the normalized Alexander polynomial of trivalent framed MOY graph, as it is directly related to the topic of MOY calculus that interests us the most. Here is our first main result.

Theorem 1.1 *For a framed trivalent MOY graph \mathbb{G} , we can fix a normalization of the Alexander polynomial $\Delta_{(\mathbb{G},c)}(t)$ so that it is a well-defined rational function invariant under framed ambient isotopy.*

While our invariant is similar to some existing invariants, which we will discuss presently, we have the following three contributions: (1) The state sum formula is new; (2) We prove the topological invariance combinatorially; namely, we prove its invariance under Reidemeister moves by applying the state sum formula; (3) The normalization of the Alexander polynomial in the framed case is new.

Our second main result is in Sect. 4. It is inspired by Murakami-Ohtsuki-Yamada’s relations in [11], where they provided a graphical definition for the $U_q(\mathfrak{sl}_n)$ -polynomial

invariants of a link for all $n \geq 2$, which is called MOY calculus. We summarize Sect. 4 as follows.

Theorem 1.2 *For a framed trivalent MOY graph \mathbb{G} , where each edge is colored by a positive integer, the invariant $\Delta_{(\mathbb{G},c)}(t)$ satisfies a series of relations, which we call MOY-type relations; conversely, these relations determine $\Delta_{(\mathbb{G},c)}(t)$. Using them we provide a graphical definition for the Alexander polynomial of a link.*

We remark that when $n = 0$, the original MOY relations in [11] gives a trivial invariant that is identically 0 for all graphs, instead of the desired Alexander invariant. There are various solutions to this issue, for example, by cutting open the graph diagrams. Nevertheless, the relations in our theorem seem to be the ones that resemble most closely the original constructions.

1.2 Relation with existing invariants

The Alexander polynomial of a knot was first defined and studied by Alexander in 1920s. Now we know there are many different interpretations coming from different backgrounds for the same invariant. Unlike the case of a knot, there exist diverse versions of Alexander polynomials for a spatial graph, which are usually distinct as invariants.

For clarity, we summarize the relation of $\langle D, c \rangle$, $\Delta_{(G,c)}(t)$ and $\Delta_{(\mathbb{G},c)}(t)$ with some existing invariants of spatial graphs.

- The state sum $\langle D, c \rangle$ simplifies the state sum proposed by the first named author in [1]. We provide an explanation in Sect. 2.
- The state sum $\langle D, c \rangle$ is also an elaboration of the Alexander polynomial defined in [5], as both the state sum in [1] and the present paper and the Alexander polynomial in [5] are special cases of the torsion invariant defined in Friedl-Juhász-Rasmussen [3].
- As an application of Theorem 1.2, the first named author showed in a subsequent paper [2, Theorem 3.2] that the Alexander polynomial $\Delta_{(\mathbb{G},c)}(t)$ is equivalent to Viro's $\mathfrak{gl}(1|1)$ -Alexander polynomial of a framed trivalent graph defined in [13], in the case that the graph has no sink or source vertices and each edge is colored by a positive integer. The key observation is that Viro's $\mathfrak{gl}(1|1)$ -Alexander polynomial also satisfies an adapted version of the MOY-type relations, which totally characterize the invariant.
- For the θ_n -graph, Litherland defined an invariant similar to $\langle D, c \rangle$ in [9].

Remark 1.3 For other versions of Alexander polynomial of spatial graphs, e.g, Kinoshita [7], as long as the definitions depend solely on the complement of the graph in S^3 , they are not equivalent to the one considered in this paper.

1.3 Organization of the paper

In Sect. 2, we introduce the preliminaries of graphs, diagrams and Kauffman states, based on which we define a state sum $\langle D, c \rangle$. We then compare it with a different state

sum proposed in [1]. The proof of their equivalence is rather lengthy, technical and isolated from the remainder of the paper, so we postpone it to the Appendix.

In Sect. 3, we prove the invariance of $\langle D, c \rangle$ under various Reidemeister moves for graph diagrams. Then we give the definition of $\Delta_{(G,c)}(t)$ and exhibit its topological invariance. Furthermore, we prove Theorem 1.1.

In Sect. 4, we elaborate on Theorem 1.2. More precisely, we wrote down a set of 10 MOY-type relations in Theorem 4.1 and prove that they uniquely characterize $\Delta_{(G,c)}(t)$ in Theorem 4.4.

In Sect. 5, we discuss some properties and topological applications of our Alexander polynomial. In particular, we give in Theorem 5.3 a necessary condition for the planarity of a graph in terms of the coefficient of its Alexander polynomial.

2 The Kauffman state sum $\langle D, c \rangle$

2.1 MOY graphs

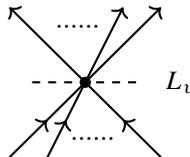
We recall the definition of an MOY graph. Our exposition and terminology is slightly different from the one used in [14].

Definition 2.1 (i) An *abstract MOY graph* is an oriented graph that equipped with a *balanced coloring* $c : E \rightarrow \mathbb{Z}_{\geq 0}$ such that for each vertex v ,

$$\sum_{e: \text{pointing into } v} c(e) = \sum_{e: \text{pointing out of } v} c(e).$$

We allow closed loops without vertices as our abstract graphs.

- (ii) An *MOY graph diagram* in \mathbb{R}^2 is an immersion of an abstract MOY graph into \mathbb{R}^2 , with crossing information and a *transverse orientation*: through each vertex v , there is a straight line L_v that separates the edges entering v and the edges leaving v .
- (iii) Two MOY graph diagrams are said to be *equivalent* if they can be connected by a finite sequence of Reidemeister moves in Fig. 2 and isotopies in \mathbb{R}^2 .
- (iv) An *MOY graph* (or a *knotted MOY graph* in Wu’s terminology in [14]) is an equivalence class of MOY graph diagrams.



Remark 2.2 An MOY graph can be regarded as a transverse spatial graph in S^3 with a balanced coloring. For the definition of a transverse spatial graph, see [5].

A vertex of valence 1 in a graph is called an *end point*; a vertex of valence greater than 1 is called an *internal vertex*. A graph is *closed* if it has no end points. A balanced

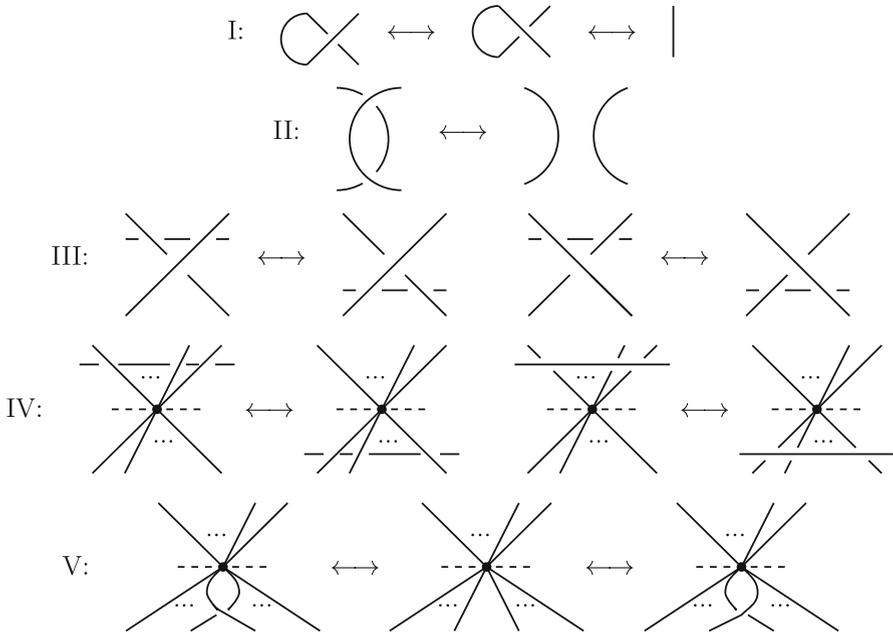


Fig. 2 Reidemeister moves for MOY graph diagrams. Suppressed orientations of the edges can be added in all compatible ways

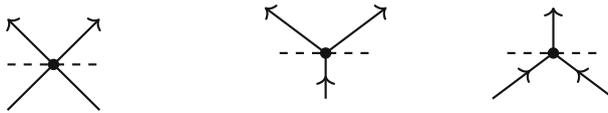


Fig. 3 A singular crossing (left) and trivalent vertices (right)

coloring is called *non-trivial* if it is not identically 0, i.e., $c(e) > 0$ for some $e \in E$; a balanced coloring is called *positive* if each edge is colored by a positive integer, i.e., $c(e) > 0$ for all $e \in E$. All graphs considered in the present paper are assumed to be closed and equipped with a non-trivial coloring. a graph is *trivalent* if all of its internal vertices have valence 3. Here are two important classes of MOY graphs.

Example 2.3 (i) Each singular knot or link gives rise to an MOY graph with a vertex corresponding to each singular crossing and a balanced coloring $c(e) = 1$ for every edge e (Fig. 3, left).

(ii) An oriented trivalent graph without source or sink has an induced transverse orientation (Fig. 3, right). With a balanced coloring, this gives an MOY graph.

2.2 Kauffman states

Let D be a graph diagram of a given MOY graph G . To define the state sum $\langle D, c \rangle$, we construct the set of states on D , which we call Kauffman states.

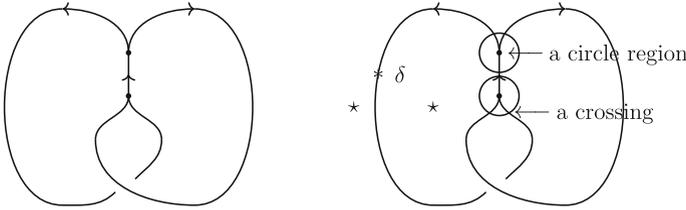


Fig. 4 Left: an MOY graph diagram D . Right: a decorated diagram (D, δ)

Definition 2.4 Starting from an MOY graph diagram D , we can obtain a *decorated diagram* (D, δ) by putting a base point δ on one of the edges in D and drawing a circle around each vertex of D . See Example 2.5.

- (i) $\text{Cr}(D)$: denotes the set of crossings, including the types $\nearrow \nwarrow$ and $\nwarrow \nearrow$ which are the double points of the diagram and the type \bigcirc which are the intersection points around each vertex between the incoming edges with the circle.
- (ii) $\text{Re}(D)$: denotes the set of regions, including the *regular regions* of \mathbb{R}^2 separated by D and the *circle regions* around the vertices. Note that there is exactly one circle region around each vertex. *Marked regions* are the regions adjacent to the base point δ , and the others are called *unmarked regions*.
- (iii) **Corners**: For a crossing of type $\nearrow \nwarrow$ or $\nwarrow \nearrow$, there are four corners around it, and we call them the *north*, *south*, *west*, and *east corners* of the crossing. Around a crossing of type \bigcirc there are three corners, and we call the one inside the circle region the *north corner*, the one on the left of the crossing the *west corner* and the one on the right the *east corner*. Note also that every corner belongs to a unique region in $\text{Re}(D)$.



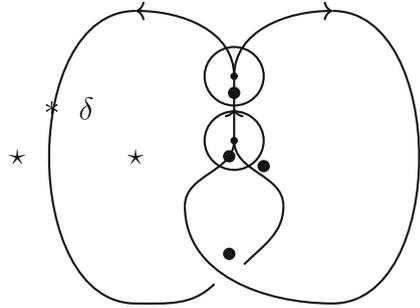
Example 2.5 In Fig. 4, we exhibit an MOY graph and a decorated diagram of it. We have $|\text{Cr}(D)| = 4$: one of which is of the type $\nearrow \nwarrow$, and three of which are of the types \bigcirc ; $|\text{Re}(D)| = 6$: four of which are regular regions, and two of which are circle regions. The base point δ is adjacent to two regions that we mark out by \star .

Calculating the Euler characteristic of \mathbb{R}^2 using D gives the following lemma.

Lemma 2.6 $|\text{Re}(D)| = |\text{Cr}(D)| + 2$ if and only if D is a connected graph diagram.

A generic base point δ is adjacent to at most two distinct regions, which will be denoted by R_u and R_v . It is possible that R_u and R_v are the same regions.

Fig. 5 A Kauffman state associated to (D, δ) is marked out by \bullet 's



Definition 2.7 A *Kauffman state*, or simply, a *state* for a decorated diagram (D, δ) is a bijective map

$$s : \text{Cr}(D) \rightarrow \text{Re}(D) \setminus \{R_u, R_v\},$$

which sends a crossing in $\text{Cr}(D)$ to one of its corners. Let $S(D, \delta)$ denote the set of all states.

Remark 2.8 It is possible that $S(D, \delta)$ is an empty set. This, for example, occurs when D is not connected or the base point δ is only adjacent to a single region.

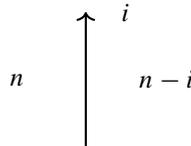
Example 2.9 The graph in Fig. 4 has three Kauffman states, one of which is illustrated in Fig. 5 by putting \bullet 's at the chosen corners.

2.3 Kauffman state sum

Suppose (D, δ) is a connected decorated diagram. Before introducing the state sum, we define the index of a regular region.

For each regular region of (D, δ) , the *index* of it is defined by the following rules.

- (i) The index of the unbounded region is set to be 0.
- (ii) The indices of the other regular regions are inductively determined by the rule as exhibited below: traversing an edge with color i in its orientation, let the difference of the index of its left-hand side region and that of its right-hand side region be i .



The definition of a balanced coloring ensures that the above rules give rise to a well-defined index for each regular region.

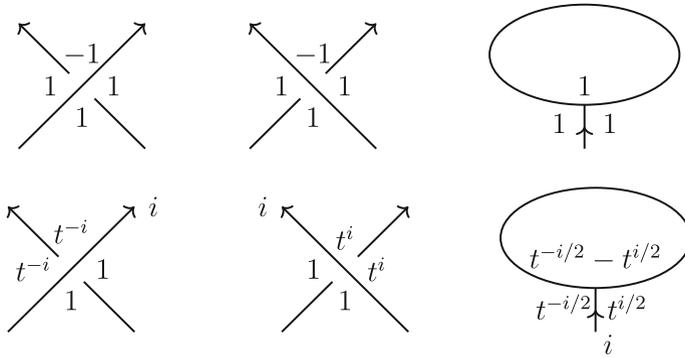


Fig. 6 The local contributions $M_{C_p}^\Delta$ (top) and $A_{C_p}^\Delta$ (bottom). Here i indicates the color of the corresponding edge

Definition 2.10 Suppose δ is on an edge with non-zero color i , and the indices of the regions adjacent to δ are n and $n - i$. Define

$$|\delta| = t^{n-i} - t^n.$$

Definition 2.11 Suppose (D, δ) is a connected decorated diagram with N crossings C_1, C_2, \dots, C_N in $\text{Cr}(D)$ and $N + 2$ regions R_1, R_2, \dots, R_{N+2} in $\text{Re}(D)$. Let c be a non-trivial coloring on D . Suppose δ is on an edge with non-zero color i . The Kauffman state sum $\langle D, c \rangle$ is defined by the following recipe.

- (i) Define the local contributions $M_{C_p}^\Delta$ and $A_{C_p}^\Delta$ associated to each corner Δ around the crossing C_p as in Fig. 6. Here $M_{C_p}^\Delta$ does not depend on the coloring, while $A_{C_p}^\Delta$ does.
- (ii) For each state $s \in S(D, \delta)$, let

$$M(s) := \prod_{p=1}^N M_{C_p}^{s(C_p)},$$

$$A(s) := \prod_{p=1}^N A_{C_p}^{s(C_p)}.$$

- (ii) The *Kauffman state sum* is defined as

$$\langle D, c \rangle := |\delta|^{-1} \sum_{s \in S(D, \delta)} M(s) \cdot A(s). \tag{1}$$

For all the cases that $S(D, \delta) = \emptyset$, we let $\langle D, c \rangle = 0$.

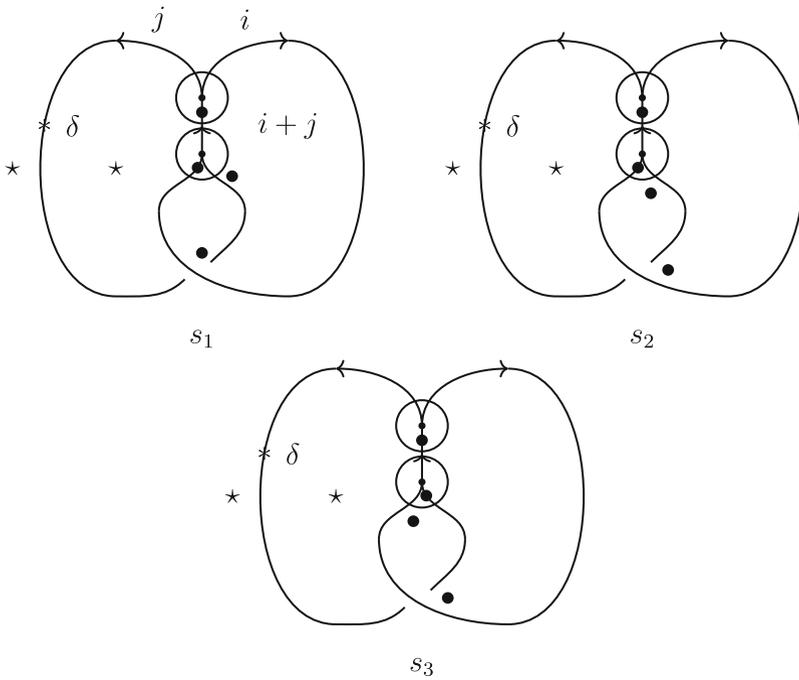


Fig. 7 The three states for the decorated diagram

Remark 2.12 While a singular knot can be regarded as a special MOY graph, the local contribution defined in Ozsváth-Stipsicz-Szabó [12, Figure 6] for the Alexander polynomial of a singular knot is different from the aforementioned construction in Fig. 6. The requirement of invariance under Reidemeister moves for MOY graphs makes it difficult to adopt a direct generalization from the former.

Example 2.13 Continuing with our previous example, we find a total of three states s_1, s_2, s_3 and represent them by \bullet 's in Fig. 7. Denote $\{k\} = t^{k/2} - t^{-k/2}$. Table 1 summarizes the computation of each state using Fig. 6.

Since $|\delta| = 1 - t^j$, we have

$$\begin{aligned}
 \langle D, c \rangle &= (1 - t^j)^{-1} [-\{i + j\}\{i\}t^{i+j/2} + \{i + j\}\{i\}t^{i-j/2} + \{i + j\}\{j\}t^{3i/2}] \\
 &= \frac{\{i + j\}[-\{i\}\{j\}t^i + \{j\}t^{3i/2}]}{-\{j\}t^{j/2}} \\
 &= -\{i + j\}t^{(i-j)/2}.
 \end{aligned}$$

2.4 Relation with the Alexander matrix

Suppose (D, δ) is a connected decorated diagram with N crossings C_1, C_2, \dots, C_N in $\text{Cr}(D)$ and $N + 2$ regions R_1, R_2, \dots, R_{N+2} in $\text{Re}(D)$. Let c be the coloring on D .

Table 1 All Kauffman states and the corresponding $M(s)$ and $A(s)$

State	$M(s)$	$A(s)$	$M(s)A(s)$
s_1	-1	$\{i + j\}i t^{j/2} t^i$	$-\{i + j\}i t^{i+j/2}$
s_2	1	$\{i + j\}i t^{-j/2} t^i$	$\{i + j\}i t^{i-j/2}$
s_3	1	$\{i + j\}j t^{i/2} t^i$	$\{i + j\}j t^{3i/2}$

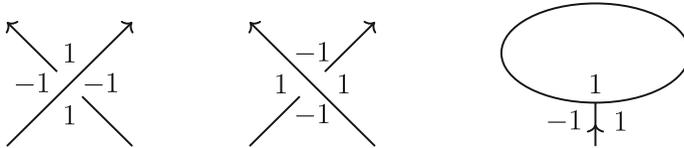
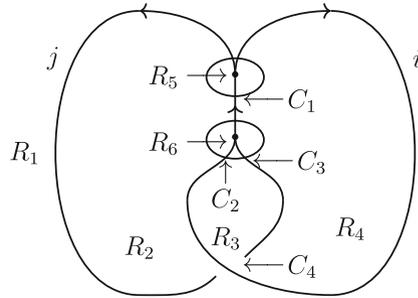


Fig. 8 The local contribution $m_{C_p}^\Delta$, where C_p is a crossing and Δ is a corner around C_p

Fig. 9 Calculating the Alexander matrix with the given labeling of the regions and crossings



Definition 2.14 The *Alexander matrix* is an $N \times (N + 2)$ matrix, denoted $A(D, c)$, whose i -th row corresponds to the crossing C_i and the j -th column corresponds to the region R_j . The (i, j) -entry of $A(D, c)$ is the sum of $m_{C_i}^\Delta A_{C_i}^\Delta$ for Δ belonging to R_j , where $m_{C_i}^\Delta$ is a local contribution specified in Fig. 8.

Example 2.15 Continuing with the previous example, we label the 4 crossings and 6 regions as in Fig. 9. We see that R_5 and R_6 are circle regions, and C_1, C_2, C_3 are crossings of type \uparrow while C_4 is a \nearrow crossing.

Now we calculate the Alexander matrix A . By definition, if C_i and R_j are not adjacent to each other, the (i, j) -entry a_{ij} of A is zero. Therefore $a_{11} = a_{13} = a_{16} = a_{21} = a_{24} = a_{25} = a_{31} = a_{32} = a_{35} = a_{45} = a_{46} = 0$.

Since the crossing C_1 meets R_2 only at its west corner, we have $a_{12} = m_{C_1}^W A_{C_1}^W = -1 \cdot t^{-(i+j)/2}$. Similarly we can calculate the other entries. The matrix is

$$A = \begin{pmatrix} 0 & -t^{-(i+j)/2} & 0 & t^{(i+j)/2} & t^{-(i+j)/2} & -t^{(i+j)/2} & 0 \\ 0 & -t^{-i/2} & t^{i/2} & 0 & 0 & 0 & t^{-i/2} - t^{i/2} \\ 0 & 0 & -t^{-j/2} & t^{j/2} & 0 & 0 & t^{-j/2} - t^{j/2} \\ -1 & 1 & -t^i & t^i & 0 & 0 & 0 \end{pmatrix}.$$

Suppose the base point δ is adjacent to two regions R_u and R_v . Let $A(D, c)\setminus(u, v)$ denote the square matrix obtained from the Alexander matrix $A(D, c)$ by removing the u -th and v -th columns, which correspond to R_u and R_v , respectively. Calculating the determinant of $A(D, c)\setminus(u, v)$ we see that

$$\det A(D, c)\setminus(u, v) = \sum_{s \in S(D, \delta)} \text{sign}(s) \cdot m(s) \cdot A(s), \tag{2}$$

where $S(D, \delta)$ is the same set of Kauffman states as before, and $\text{sign}(s)$ is the sign of the state s as a permutation with respect to the given orders of the elements in $\text{Cr}(D)$ and $\text{Re}(D)\setminus\{R_u, R_v\}$. The relation of the old state sum (2) and $\langle D, c \rangle$ is as follows.

Proposition 2.16 (Appendix) *For any $s \in S(D, \delta)$, the value*

$$\frac{M(s)}{\text{sign}(s) \cdot m(s)} \in \{1, -1\}$$

does not depend on the choice of the state s . As a result,

$$\langle D, c \rangle = \pm |\delta|^{-1} \cdot \det A(D, c)\setminus(u, v),$$

where the sign \pm only depends on the ordering in $\text{Cr}(D)$ and $\text{Re}(D)\setminus\{R_u, R_v\}$.

Remark 2.17 Besides the above combinatorial definition of $A(D, c)$, the Alexander matrix also has the following topological description. Let G be the spatial graph in S^3 determined by the graph diagram D . Denote X the complement of G in S^3 . Using the coloring c we can construct a map

$$\psi_c : \pi_1(X) \rightarrow \mathbb{Z},$$

where the meridian of each edge is sent to the color of the edge. One can construct an abelian covering space $p : \tilde{X} \rightarrow X$ from ψ_c . Then, $A(D, c)$ is a presentation matrix for the module $H_1(\tilde{X}, p^{-1}(\partial_{\text{in}}(X)))$, where $\partial_{\text{in}}(X) \subset \partial X$ is a certain submanifold associated with the transverse orientation of D and δ . See [1] for more details, where the state sum (2) was first proposed.

2.5 The choice of the base point

In this section, we prove the independence of $\langle D, c \rangle$ on the choice of the base point δ .

Proposition 2.18 *The state sum $\langle D, c \rangle$ does not depend on the choice of the location of δ among the edges with non-zero colors.*

Proof The proof is adapted from [6, Theorem 3.3]. Suppose δ and δ' are two different choices of base points. Suppose that R_u, R_v are the two regions adjacent to δ and R_u is on the left-hand side when traversing the edge containing δ in its orientation, and that R_w, R_z are the two regions adjacent to δ' . Let $\langle D, c \rangle_\delta$ be the state sum calculated by using δ . Our goal is to show that $\langle D, c \rangle_\delta = \langle D, c \rangle_{\delta'}$. The strategy is as follows. We first prove that $\langle D, c \rangle_\delta = \pm \langle D, c \rangle_{\delta'}$ using the Alexander matrix in Definition 2.14. Then we remove the ambiguity of the sign.

Without loss of generality, we assume that R_1, R_2, \dots, R_M are the regular regions and $R_{M+1}, R_{M+2}, \dots, R_{N+2}$ are the circle regions in $\text{Re}(D)$. Consider the Alexander matrix $A = A(D, c)$. Let A_q be the q -th column corresponding to the region R_q . We claim that $\sum_{q=1}^{N+2} A_q = 0$ since for each crossing C_p the sum of $m_{C_p}^\Delta A_{C_p}^\Delta$ over all corners of C_p is zero. From this identity we have

$$\det A \setminus (u, v) = (-1)^{w-v} \det A \setminus (u, w) - (-1)^{w-u} \det A \setminus (v, w) \tag{*}$$

for $1 \leq u < v < w \leq M$.

On the other hand, we have $\sum_{q=1}^M t^{\text{ind}(R_q)} A_q = 0$. Therefore

$$\sum_{q=1}^M t^{\text{ind}(R_q) - \text{ind}(R_u)} A_q = 0$$

for any fixed $1 \leq u \leq M$.

As a result, for any $1 \leq u, w \leq M$, we have

$$-t^{\text{ind}(R_w) - \text{ind}(R_u)} A_w = \sum_{\substack{1 \leq q \leq M \\ q \neq w}} t^{\text{ind}(R_q) - \text{ind}(R_u)} A_q.$$

Bringing the scalar $-t^{\text{ind}(R_w) - \text{ind}(R_u)}$ to the column A_w in $A \setminus (u, v)$, we get

$$\begin{aligned} & -t^{\text{ind}(R_w) - \text{ind}(R_u)} \det A \setminus (u, v) \\ &= (-1)^{w-u} \det A \setminus (v, w) - t^{\text{ind}(R_v) - \text{ind}(R_u)} (-1)^{w-v} \det A \setminus (u, w) \end{aligned}$$

for $1 \leq u < v < w \leq M$. This together with (*) implies that

$$(t^{\text{ind}(R_w)} - t^{\text{ind}(R_v)}) \det A \setminus (u, v) = (-1)^{w-u} (t^{\text{ind}(R_v)} - t^{\text{ind}(R_u)}) \det A \setminus (v, w).$$

Applying the above relation twice, we end up with

$$(t^{\text{ind}(R_z)} - t^{\text{ind}(R_w)}) \det A \setminus (u, v) = (-1)^{(w+z)-(u+v)} (t^{\text{ind}(R_v)} - t^{\text{ind}(R_u)}) \det A \setminus (w, z),$$

for any $1 \leq u < v \leq w < z \leq M$.

By definition we have $|\delta| = t^{\text{ind}(R_v)} - t^{\text{ind}(R_u)}$ and $|\delta'| = t^{\text{ind}(R_z)} - t^{\text{ind}(R_w)}$. Choose an ordering of crossings in $\text{Cr}(D)$ so that $\det A \setminus (u, v) = |\delta| \langle D, c \rangle_\delta$. Then we have $\det A \setminus (w, z) = (-1)^\epsilon |\delta'| \langle D, c \rangle_{\delta'}$ for $\epsilon = 0$ or 1 . The sign ϵ depends only on the ordering of the regions and crossings. In summary, we have proved that

$$\langle D, c \rangle_\delta = (-1)^{(w+z)-(u+v)+\epsilon} \langle D, c \rangle_{\delta'} = \pm \langle D, c \rangle_{\delta'}. \tag{**}$$

Next, we remove “ \pm ” from (**). Let us first consider the case that D is a plane diagram. As $\langle D, c \rangle_\delta$ is zero if and only if $\langle D, c \rangle_{\delta'}$ is zero, we assume that $\langle D, c \rangle_\delta$ is non-zero. Note that there are only crossing of type \bigcirc for a plane diagram. From Formula (1) of $\langle D, c \rangle$, we see that the coefficient of the lowest degree term of $|\delta| \langle D, c \rangle_\delta$ is positive. Since $\text{ind}(R_u) > \text{ind}(R_v)$ and $\text{ind}(R_w) > \text{ind}(R_z)$, the coefficients of the lowest degree terms of both $|\delta'| |\delta| \langle D, c \rangle_\delta$ and $|\delta| |\delta'| \langle D, c \rangle_{\delta'}$ are also positive. This together with (**) implies that

$$\langle D, c \rangle_\delta = \langle D, c \rangle_{\delta'}.$$

Hence, the location of δ does not affect the value of $\langle D, c \rangle$ when D is a plane MOY diagram.

For an arbitrary MOY graph diagram that possibly has crossings of type  or , we connect it with plane graphs by the following skein-type relations.

$$\begin{aligned} \left\langle \begin{array}{c} j \quad i \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} \right\rangle &= \frac{-t^{-\frac{i+j}{2}}}{\{i\}\{j\}} \cdot \left\langle \begin{array}{c} j \quad i \\ \uparrow \quad \uparrow \\ \diagdown \quad \diagup \\ \uparrow \quad \uparrow \\ i \quad j \end{array} \right\rangle - \frac{t^{-\frac{j}{2}}}{\{i\}} \cdot \left\langle \begin{array}{c} j \quad i \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ i \quad j \end{array} \right\rangle, \\ \left\langle \begin{array}{c} j \quad i \\ \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right\rangle &= \frac{-t^{\frac{i+j}{2}}}{\{i\}\{j\}} \cdot \left\langle \begin{array}{c} j \quad i \\ \uparrow \quad \uparrow \\ \diagdown \quad \diagup \\ \uparrow \quad \uparrow \\ i \quad j \end{array} \right\rangle - \frac{t^{\frac{j}{2}}}{\{i\}} \cdot \left\langle \begin{array}{c} j \quad i \\ \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ i \quad j \end{array} \right\rangle, \end{aligned}$$

where $j \geq i \geq 0$ and $\{i\} = t^{i/2} - t^{-i/2}$. For $j \leq i$ we can also establish similar relations. Such kinds of relations will be discussed in great details in Sect. 4, so we omit their proofs here. Note that all the coefficients appearing on the right-hand side of the equations do not depend on δ , and $|\delta|$ is identical for all the graphs showing here. As we just proved that $\langle D, c \rangle$ is independent of the choice of δ in the case of plane graphs, we conclude that the same is true for general MOY graph diagrams from the relations above. \square

3 The Alexander polynomial

In this section, we study the invariance of the Kauffman state sum under the Reidemeister moves. This enables us to define an Alexander polynomial $\Delta_{(G,c)}(t)$ for all MOY graphs. Then we provide a normalization $\Delta_{(\mathbb{G},c)}(t)$ for a framed trivalent graph \mathbb{G} .

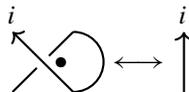
3.1 The invariance of the Kauffman state sum $\langle D, c \rangle$

Proposition 3.1 *The Kauffman state sum $\langle D, c \rangle$ is invariant under the Reidemeister moves (II) – (V) in Fig. 2, and its variations under Reidemeister move (I) are given as below.*

$$t^i \langle \text{Diagram 1} \rangle = \langle \text{Diagram 2} \rangle = \langle \text{Diagram 3} \rangle = t^{-i} \langle \text{Diagram 4} \rangle = \langle \text{Diagram 5} \rangle.$$

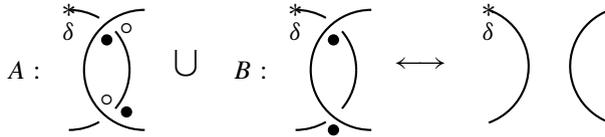
Proof The main trick is to apply Proposition 2.18 and choose the location of the base point δ that is most convenient for calculations.

For move (I), we assume that it occurs away from the base point. Take the following figure for example, where the left-hand diagram is D and the right-hand one is D' . There is a one-to-one correspondence between the states of D and those of D' since the additional region of D has only one corner. It then follows from (1) and the local contributions described in Fig. 6 that $\langle D, c \rangle = t^i \langle D', c \rangle$. The proofs for the other cases of move (I) follow from similar arguments.

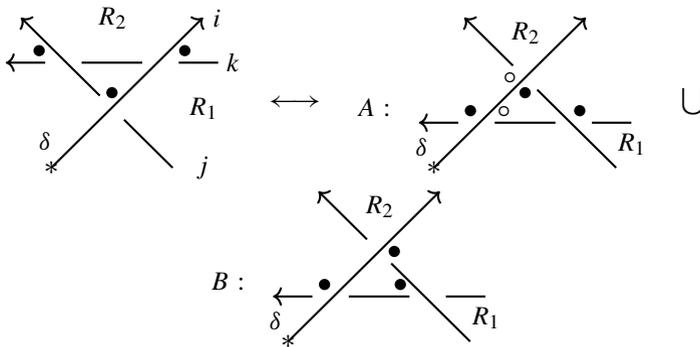


The proof of the invariance of $\langle D, c \rangle$ under moves (II), (III) is similar to the proof of [6, Theorem 4.3]. For move (II), we place the base point in a position as below. Suppose the left-hand graph is D and the right-hand one is D' . We split the set $S(D, \delta)$ as a disjoint union of A and B , where B consists of all the states which assign the bottom crossing of D to its south corner and A is the complement of B . The contribution of A to $\langle D, c \rangle$ is 0 since the states are paired up (white states and black states) and the contribution from each pair adds up to 0. There is a one-to-one correspondence

between B and $S(D', \delta)$. Since for each state in B the contribution from the two crossings involved is 1, we conclude that $\langle D, c \rangle = \langle D', c \rangle$.



For move (III), we consider the move in the following figure for illustration, and the invariance under other cases can be proved similarly. We assume that δ is located on the edge with color i . Around the triangle where the move occurs, we have three crossings and five local regions not adjacent to δ (some of them can be adjacent to δ or belong to the same region globally), so for each state there are exactly two blank local regions in the sense that there is no crossing assigned to them around the triangle. Suppose the left-hand diagram is D and the right-hand one is D' . We compare the states of D and D' that have the same blank local regions around the triangle and are identical outside the triangle.



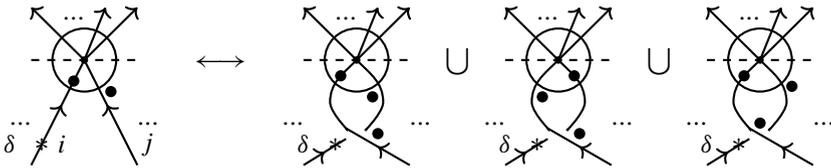
For instance, consider the states whose blank regions are R_1 and R_2 . Each of such states of D has a unique local assignment as in the left-hand figure. One can check that the contribution from the three crossings is $t^{-(i+j)}$. On D' there are three types of states with the indicated blank regions. The set A comprises the first two types of states. Their contribution to $\langle D', c \rangle$ vanishes since the states are paired up (white states and black states) and the contribution from each pair is zero. For each state of B , the contribution from the three crossings is $t^{-(i+j)}$. As a result, for the states with the blank regions R_1 and R_2 , their contributions to the state sum on both sides are the same. For the states with different blank regions, we can do similar calculations to verify that the contributions from both sides are identical as well. Consequently, we have $\langle D, c \rangle = \langle D', c \rangle$.

For move (V), the twist creates a crossing on the diagram. If the two edges of the crossing are outgoing edges, we assume that δ is away from the twist. Since the newly created crossing must be assigned to its south corner whose contribution is 1, we see that $\langle D, c \rangle$ is invariant. Now we consider the case that the two edges of the crossing are incoming edges. We assume that δ is located as in the figure below. Suppose the left-hand diagram is D and the right-hand one is D' . The assignment of the two crossings

of type $\textcircled{1}$ on D is unique, and their contribution to $\langle D, c \rangle$ is $t^{j/2}(t^{-i/2} - t^{i/2})$. On D' there are three types of local assignments, and the sum of their contributions to $\langle D', c \rangle$ is

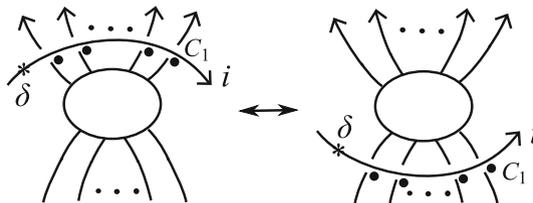
$$t^{-i/2+j}(t^{-j/2} - t^{j/2}) + t^{j/2+j}(t^{-i/2} - t^{i/2}) - t^{i/2+j}(t^{-j/2} - t^{j/2}) = t^{j/2}(t^{-i/2} - t^{i/2}).$$

Therefore we see that $\langle D, c \rangle = \langle D', c \rangle$.

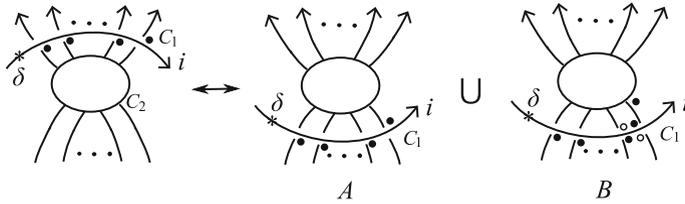


For the invariance under move (IV), we again argue by comparing the states of two diagrams before and after the move. The proof is tedious nonetheless straightforward, so we will only write down the idea but omit the details. Take the following move for example, where the transverse edge passes over other edges. We assume that δ is on the transverse edge with color i as indicated. Consider the crossings of the left-hand diagram D lying on the transverse edge. There are essentially only two different assignments of them that we need to consider, the first of which assigns the crossing C_1 to its east corner, and the second to its north corner. Other assignments, if any, always appear in pairs and the contribution of each pair to $\langle D, c \rangle$ vanishes.

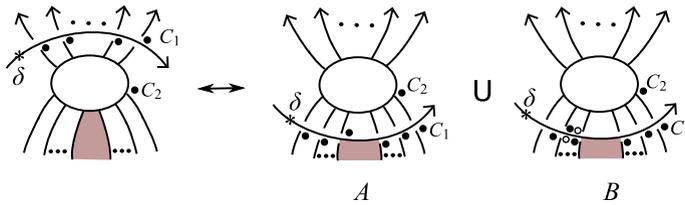
Consider the states where C_1 is assigned to its east corner. Then the states of the right-hand diagram D' which have the same blank regions must assign C_1 to its east corner as well, as shown below. We can then check that their contributions to the state sum on both sides are the same.



Consider the states of D where C_1 is assigned to its north corner and C_2 is not assigned to its east corner. In this case, the states of D' with the same blank regions are divided into two sets A and B , where the each state in A assigns C_1 to its north corner and each state in B assigns C_1 to its west or south corner, as shown below. The states in B , no matter how many, always appear in pairs and the sum of their contributions to the state sum is zero. Hence, we also have in this case that the contribution from the states of D equals that from A .



Consider the states of D where C_1 is assigned to its north corner and C_2 is assigned to its east corner. For each of such states, there must be a blank region under the circle region. We shadow this region, as the example shown below. The states of D' which provide the same blank regions are divided into two sets A and B , where each state in A sends the crossing in the northwest corner of the shadow region to its south corner, and each state in B sends the crossing to its west or south corner. The states in B , no matter how many, can always be paired up and the sum of their contributions to the state sum is zero. This again shows that the contribution from the states of D in this case equals that from A . Combining the discussion above, we complete the proof of the invariance of $\langle D, c \rangle$ under move (IV).



□

Definition 3.2 Let (G, c) be an MOY graph. The Alexander polynomial of (G, c) , denoted $\Delta_{(G,c)}(t)$, can be defined from any of its graph diagram (D, c) by

$$\Delta_{(G,c)}(t) := \frac{\langle D, c \rangle}{(t^{-1/2} - t^{1/2})^{|V|-1}}, \tag{3}$$

where $|V|$ is the number of vertices of G .

The term $(t^{-1/2} - t^{1/2})^{|V|-1}$ in the denominator of Formula (3) is clearly an invariant of the graph. The advantage for introducing this factor in the definition of the Alexander polynomial will become transparent when we discuss the properties of $\Delta_{(D,c)}(t)$ for plane graphs (Theorem 5.3).

Together with Proposition 2.18 and Proposition 3.1, it implies that $\Delta_{(G,c)}(t)$ is well-defined up to a power of t , just like the classical Alexander polynomial of links.

3.2 The normalization of the Alexander polynomial

In this part, we present a way of normalizing the Alexander polynomial $\Delta_{(G,c)}(t)$ that eliminates the t^k -ambiguity in its definition when we fix a framing on G . Whereas our

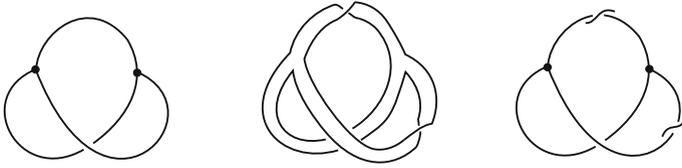


Fig. 10 A framed graph and its graph diagram

idea is applicable to all framed MOY graphs, we will focus on the class of trivalent graphs, because it is the natural setting for the topic in Sect. 4 on MOY calculus.

Recall that a *framing* of G is an embedded compact surface $F \subset S^3$ in which G is sitting as a deformation retract. A *framed graph* is a graph equipped with a framing. More precisely, each vertex of G is replaced by a disk in F where the vertex is the center of the disk, and each edge of G is replaced by a strip $[0, 1] \times [0, 1]$ where $[0, 1] \times \{0, 1\}$ is attached to the boundaries of its adjacent vertex disks and the edge is $\{\frac{1}{2}\} \times [0, 1]$. Each graph diagram of G in \mathbb{R}^2 has a *blackboard framing*, whose projection in \mathbb{R}^2 is the tubular neighborhood of the graph diagram in \mathbb{R}^2 . The difference between a generic framing and a blackboard framing can be represented on the diagram by introducing the symbols $\text{---}\swarrow\text{---}$ and $\text{---}\searrow\text{---}$, which mean a positive half twist and a negative half twist respectively at the fragment. Therefore we can use a graph diagram with $\text{---}\swarrow\text{---}$ and $\text{---}\searrow\text{---}$ to represent a framed graph, as indicated in Fig. 10.

Our idea of normalization is simple and natural. Starting from a framed graph diagram, we hope to construct an appropriate factor that cancels the $t^{\pm i}$ factor in the Alexander polynomial coming from Reidemeister move (I) in Proposition 3.1 while keeping invariant under the other types of Reidemeister moves.

For framed trivalent graphs, the following result is well-known to experts (see for example [13]).

Lemma 3.3 *Any two graph diagrams for a framed trivalent graph can be connected by a sequence of Reidemeister moves in Fig. 11.*

From now on, we use the blackboard-bold letters \mathbb{G} and \mathbb{D} to denote a framed trivalent graph and diagram, respectively. We first define a framing factor.

Definition 3.4 For each edge e with color i , let

$$\mathcal{F}(e) = t^{i/4[(\sharp\text{---}\searrow\text{---}) - (\sharp\text{---}\swarrow\text{---})]},$$

where $\sharp\text{---}\searrow\text{---}$ and $\sharp\text{---}\swarrow\text{---}$ denote the number of $\text{---}\searrow\text{---}$ and that of $\text{---}\swarrow\text{---}$ on e , respectively. The *framing factor* associated to (\mathbb{D}, c) is define as

$$\mathcal{F}(\mathbb{D}, c) = \prod_{e \in E} \mathcal{F}(e),$$

where E is the edge set of \mathbb{D} .

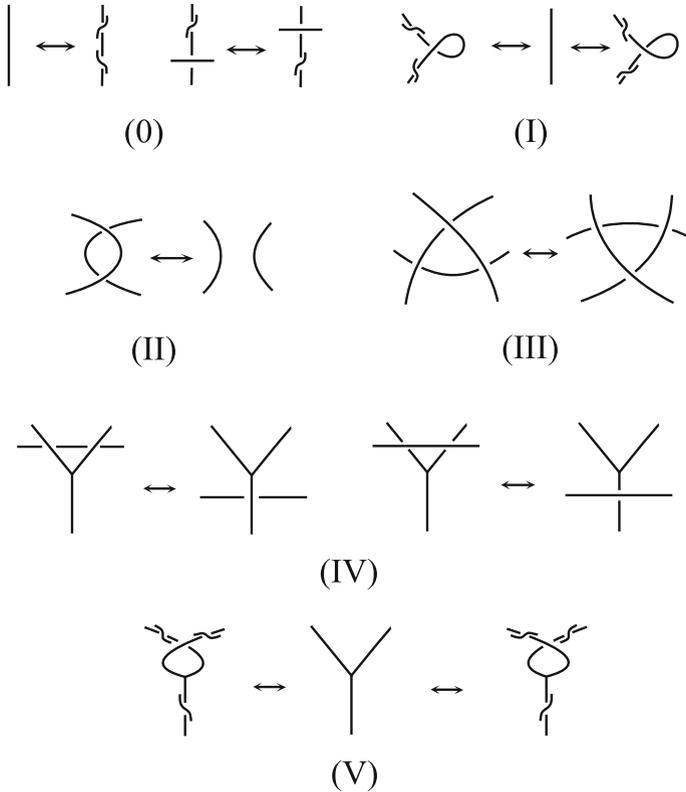


Fig. 11 The Reidemeister moves for framed trivalent graphs

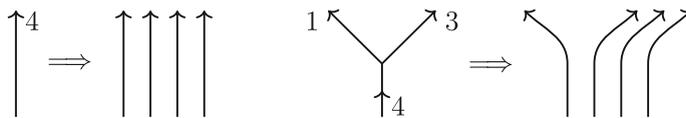


Fig. 12 The way we obtain the link $L_{(D,c)}$ from a trivalent diagram D

From D we can obtain a diagram of an oriented link $L_{(D,c)}$ by the local transformation depicted in Fig. 12, where we replace each edge of color i by i parallel strands.

Definition 3.5 We define the *curliness* of D with respect to c by

$$C(D, c) = t^{1/2[\# \text{ } \curvearrowright - (\# \text{ } \curvearrowleft)]},$$

where $\# \curvearrowright$ and $\# \curvearrowleft$ denote respectively the numbers of counter-clockwise and clockwise oriented curves that we obtain after resolving the crossings of $L_{(D,c)}$ in the oriented way.

As the notation indicates, $\mathcal{C}(D, c)$ does not depend on framing. Roughly speaking, $\mathcal{C}(D, c)$ measures the total “winding number” of the diagram D with the weight c . Our next lemma is evident from this viewpoint.

Lemma 3.6 *The curliness $\mathcal{C}(D, c)$ satisfies the following local relation:*

$$c \left(\begin{array}{c} \curvearrowright \\ i \end{array} \right) \left(\begin{array}{c} \curvearrowleft \\ i \end{array} \right) = t^{-i/2} \mathcal{C} \left(\begin{array}{c} \curvearrowright \\ i \quad \curvearrowleft \\ i \end{array} \right).$$

The curliness $\mathcal{C}(D, c)$ is invariant under moves (II) and (III). Its changes under move (I) are given as follows.

$$\begin{aligned} t^{-i/2} \mathcal{C} \left(\begin{array}{c} \curvearrowright \\ i \end{array} \right) &= t^{-i/2} \mathcal{C} \left(\begin{array}{c} \curvearrowright \\ i \end{array} \right) = t^{i/2} \mathcal{C} \left(\begin{array}{c} \curvearrowleft \\ i \end{array} \right) = t^{i/2} \mathcal{C} \left(\begin{array}{c} \curvearrowleft \\ i \end{array} \right) \\ &= \mathcal{C} \left(\begin{array}{c} \uparrow \\ i \end{array} \right). \end{aligned}$$

Proof One can check the relations by direct computations. □

Definition 3.7 For a framed trivalent MOY graph diagram (\mathbb{D}, c) , define the *normalized Alexander polynomial* by

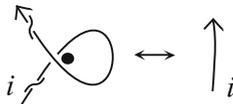
$$\Delta_{(\mathbb{D}, c)}(t) := \frac{\langle D, c \rangle}{(t^{-1/2} - t^{1/2})^{|V|-1}} \cdot \mathcal{F}(\mathbb{D}, c) \mathcal{C}(D, c). \tag{4}$$

Now we prove Theorem 1.1.

Theorem 3.8 (Theorem 1.1) *The function $\Delta_{(\mathbb{D}, c)}(t)$ is invariant under the Reidemeister moves in Fig. 11. Therefore, it is a topological invariant of the corresponding framed trivalent graph (\mathbb{G}, c) .*

Proof It is straightforward to check that $\Delta_{(\mathbb{D}, c)}(t)$ is invariant under move (0).

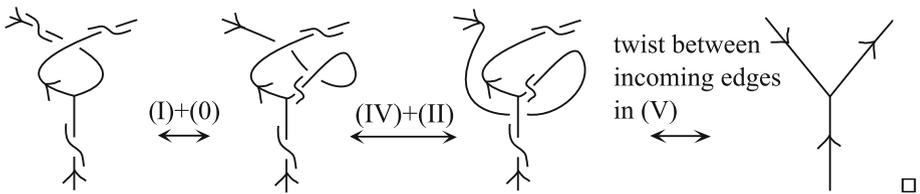
For move (I), we prove the move in the following figure for illustration, where \mathbb{D} denotes the left-hand side diagram and \mathbb{D}' denotes the right-hand side diagram. One can check that $\mathcal{F}(\mathbb{D}, c) = t^{-i/2} \mathcal{F}(\mathbb{D}', c)$ and $\mathcal{C}(D, c) = t^{-i/2} \mathcal{C}(D', c)$, while Proposition 3.1 implies that the state sums satisfy $\langle D, c \rangle = t^i \langle D', c \rangle$. Therefore $\Delta_{(\mathbb{D}, c)}(t) = \Delta_{(\mathbb{D}', c)}(t)$. The proof for the other cases of move (I) are similar. However, note that for the other cases of move (I), $\mathcal{F}(\mathbb{D}, c)$ and $\mathcal{C}(D, c)$ do not necessarily change in the same way.



From Lemma 3.6, we see that $\mathcal{F}(\mathbb{D}, c)$ and $\mathcal{C}(D, c)$ are invariant under moves (II) (III) and (IV). Proposition 3.1 implies the same fact for $\langle D, c \rangle$. Therefore, $\Delta_{(\mathbb{D}, c)}(t)$ is also invariant under moves (II) (III) and (IV).

For move (V), we divide the discussion into two cases depending on the orientation of the edges. The first case is when the two edges where the twist occurs point both inward or outward. This belongs to move (V) in Fig. 2. Proposition 3.1 implies that $\langle D, c \rangle$ is invariant, and we can check from the definition that $\mathcal{F}(\mathbb{D}, c)$ and $\mathcal{C}(D, c)$ are also invariant. As a result, $\Delta_{(\mathbb{D}, c)}(t)$ is invariant.

The second case is when one edge around the twist points inward and the other edge points outward. As this move can be realized as a composition of the moves discussed above, the invariance of $\Delta_{(\mathbb{D}, c)}$ follows. See the following figure.



Thus, we define the normalized Alexander polynomial of framed trivalent graph (\mathbb{G}, c) by

$$\Delta_{(\mathbb{G}, c)}(t) := \Delta_{(\mathbb{D}, c)}(t) = \frac{\langle D, c \rangle}{(t^{-1/2} - t^{1/2})^{|V|-1}} \cdot \mathcal{F}(\mathbb{D}, c)\mathcal{C}(D, c), \tag{5}$$

where \mathbb{D} is an arbitrary diagram that represents \mathbb{G} .

Remark 3.9 In summary, for an MOY graph G without framing, we have only an Alexander polynomial $\Delta_{(G, c)}(t)$ that is well-defined up to t^k . When we consider a framed trivalent graph \mathbb{G} , we obtain a normalized Alexander polynomial $\Delta_{(\mathbb{G}, c)}(t)$ without t^k -ambiguity. We will use both $\Delta_{(G, c)}(t)$ and $\Delta_{(\mathbb{G}, c)}(t)$ in subsequent sections for different purposes.

4 MOY-type relations

In this section, we prove Theorem 1.2. This result originates from Murakami-Ohtsuki-Yamada’s graphic calculus [11] for $U_q(\mathfrak{sl}_n)$ -polynomial invariants. Here we give a version of the theory for $\Delta_{(\mathbb{G}, c)}(t)$.

4.1 MOY-type relations.

Theorem 4.1 *Suppose \mathbb{D} is a framed trivalent graph diagram with a positive coloring c . The normalized Alexander polynomial $\Delta_{(\mathbb{D}, c)}(t)$ satisfies the following relations,*

where (\mathbb{D}) represents $\Delta_{(\mathbb{D},c)}(t)$, and $[k] = \frac{t^{k/2} - t^{-k/2}}{t^{1/2} - t^{-1/2}}$ for $k \in \mathbb{Z}$.

(i) $\left(\begin{array}{c} \text{circle with arrow} \\ i \end{array} \right) = \left(\begin{array}{c} \text{circle with arrow} \\ i \end{array} \right) = \frac{1}{[i]}$, for $i > 0$.

(ii) $(\mathbb{D}) = 0$ if \mathbb{D} is a disconnected graph.

(iii) $\left(\begin{array}{c} i \\ \text{fork} \end{array} \right) = t^{-i/4} \cdot \left(\begin{array}{c} i \\ \text{bar} \end{array} \right)$, $\left(\begin{array}{c} i \\ \text{bar} \end{array} \right) = t^{i/4} \cdot \left(\begin{array}{c} i \\ \text{fork} \end{array} \right)$.

(iv) When $i \leq j$:

$$\left(\begin{array}{c} j \quad i \\ \text{cross} \end{array} \right) = \frac{-t^{-\frac{i+j}{2}}}{[i][j]} \cdot \left(\begin{array}{c} j \quad i \\ \text{bar with arrow} \end{array} \right) + \frac{t^{-\frac{j}{2}}}{[i][i+j]} \cdot \left(\begin{array}{c} j \quad i \\ \text{merge} \end{array} \right),$$

$$\left(\begin{array}{c} j \quad i \\ \text{cross} \end{array} \right) = \frac{-t^{\frac{i+j}{2}}}{[i][j]} \cdot \left(\begin{array}{c} j \quad i \\ \text{bar with arrow} \end{array} \right) + \frac{t^{\frac{j}{2}}}{[i][i+j]} \cdot \left(\begin{array}{c} j \quad i \\ \text{merge} \end{array} \right).$$

When $j \leq i$:

$$\left(\begin{array}{c} j \quad i \\ \text{cross} \end{array} \right) = \frac{-t^{-\frac{i+j}{2}}}{[i][j]} \cdot \left(\begin{array}{c} j \quad i \\ \text{bar with arrow} \end{array} \right) + \frac{t^{-\frac{j}{2}}}{[j][i+j]} \cdot \left(\begin{array}{c} j \quad i \\ \text{merge} \end{array} \right),$$

$$\left(\begin{array}{c} j \quad i \\ \diagdown \quad \diagup \\ / \quad \diagdown \end{array} \right) = \frac{-t^{\frac{i+j}{2}}}{[i][j]} \cdot \left(\begin{array}{c} j \quad i \\ \uparrow \quad \uparrow \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ i \quad j \end{array} \right) + \frac{t^{\frac{i}{2}}}{[j][i+j]} \cdot \left(\begin{array}{c} j \quad i \\ \diagdown \quad \diagup \\ \bullet \\ \downarrow \\ i+j \\ \bullet \\ \diagup \quad \diagdown \\ i \quad j \end{array} \right).$$

$$(v) \quad \left(\begin{array}{c} j \uparrow \\ i+j \uparrow \quad i \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ j \uparrow \end{array} \right) = \left(\begin{array}{c} j \uparrow \\ i \leftarrow \quad i+j \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ j \uparrow \end{array} \right) = [j][i+j] \cdot \left(\begin{array}{c} \uparrow \\ j \end{array} \right).$$

$$(vi) \quad \left(\begin{array}{c} i \uparrow \\ i-j \quad j \\ \bullet \quad \bullet \\ \downarrow \\ i \end{array} \right) = [i]^2 \cdot \left(\begin{array}{c} \uparrow \\ i \end{array} \right), \text{ where } i \geq j.$$

$$(vii) \quad \left(\begin{array}{c} j \quad i+j \quad i \\ \uparrow \quad \uparrow \quad \uparrow \\ \bullet \quad \bullet \quad \bullet \\ \downarrow \quad \downarrow \quad \downarrow \\ i \quad j \quad i \\ \uparrow \quad \uparrow \quad \uparrow \\ j \quad i+j \end{array} \right) = \frac{[i+j]^3}{[i]} \cdot \left(\begin{array}{c} j \quad i \\ \uparrow \quad \uparrow \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ i-j \end{array} \right) + [j]^2 [i+j]^2 \cdot \left(\begin{array}{c} \uparrow \quad \downarrow \\ j \quad i \end{array} \right), \text{ where } i \geq j.$$

$$(viii) \quad \left(\begin{array}{c} i \quad j \quad k \\ \uparrow \quad \uparrow \quad \uparrow \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ i+j \quad i+j+k \end{array} \right) = \frac{[i+j]}{[j+k]} \cdot \left(\begin{array}{c} i \quad j \quad k \\ \uparrow \quad \uparrow \quad \uparrow \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ j+k \quad i+j+k \end{array} \right),$$

$$\left(\begin{array}{c} i+j \quad i+j+k \\ \uparrow \quad \uparrow \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ i \quad j \quad k \end{array} \right) = \frac{[i+j]}{[j+k]} \cdot \left(\begin{array}{c} i+j+k \\ \uparrow \quad \uparrow \\ \bullet \quad \bullet \\ \downarrow \quad \downarrow \\ i \quad j \quad k \end{array} \right).$$

$$\begin{aligned}
 (ix) \quad & \left(\begin{array}{c} i+k-l \quad j+l-k \\ \uparrow \quad \quad \uparrow \\ \text{---} \xrightarrow{l} \text{---} \\ \uparrow \quad \quad \uparrow \\ i+k \quad j-k \\ \uparrow \quad \quad \uparrow \\ i \quad \quad j \\ \text{---} \xleftarrow{k} \text{---} \\ \uparrow \quad \quad \uparrow \end{array} \right) = \frac{[j][l][i+k]}{[i+j]} \cdot \left(\begin{array}{c} i+k-l \quad j+l-k \\ \quad \quad \quad \uparrow \quad \uparrow \\ \quad \quad \quad \text{---} \quad \text{---} \\ \quad \quad \quad \uparrow \\ i+j \\ \quad \quad \quad \downarrow \quad \downarrow \\ i \quad \quad j \end{array} \right) \\
 & + [i+k][j-k] \cdot \left(\begin{array}{c} i+k-l \quad j+l-k \\ \uparrow \quad \quad \uparrow \\ \quad \quad \quad \text{---} \xleftarrow{k-l} \text{---} \\ \uparrow \quad \quad \uparrow \\ i \quad \quad j \end{array} \right), \text{ where } j \geq k \geq l. \\
 (x) \quad & \left(\begin{array}{c} i \quad j \\ | \quad | \\ \diagdown \quad \diagup \\ | \quad | \\ i \quad j \end{array} \right) = [i][j] \cdot \left(\begin{array}{c} i \quad j \\ | \quad | \\ | \quad | \\ i \quad j \end{array} \right).
 \end{aligned}$$

Proof Like other literature on MOY calculus, it is customary to treat a loop as a special vertex-less trivalent MOY graph. We can then evaluate the Kauffman state sum formula in relation (i) as follows. Note that both the set of crossings and the set of unmarked regions (after we place a base point on the circle) are empty. Hence, the set of Kauffman states $S(D)$ is a singleton set whose unique element is the unique map $s : \emptyset = \text{Cr}(D) \rightarrow \text{Re}(D) \setminus \{R_u, R_v\} = \emptyset$. Then, according to Definition 2.11, $M(s)$ and $A(s)$ are both the trivial product, i.e., 1. With $|V| = 0$, $\mathcal{F}(\mathbb{D}, c) = 1$, $\mathcal{C}(D, c) = t^{-i/2}$ and $|\delta| = t^{-i} - 1$ for the clockwise diagram and similar computation for the counter-clockwise diagram, Definition 3.7 gives the desired relation (i).

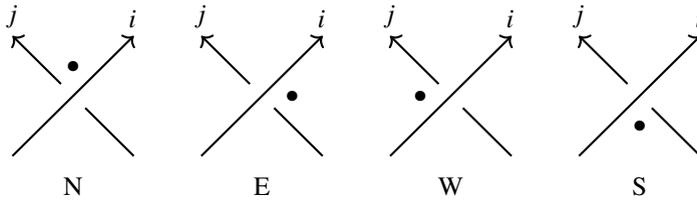
Relation (ii) follows from the definition in Section 2.2, and (iii) follows from the definition of $\mathcal{F}(\mathbb{D}, c)$.

To prove the other relations, the basic strategy is similar to Proposition 3.1 - we establish a bijection between the terms in the state sum $\langle D, c \rangle = |\delta|^{-1} \sum_{s \in S(D, \delta)} M(s) \cdot A(s)$ of the left-hand and right-hand side of each relation.

As an illustration, let us prove the first relation in (iv). Note that we used similar relations in the proof of Proposition 2.18 without giving a proof. Suppose the left-hand diagram in the first relation of (iv) is D and the right-hand two diagrams are D_1 and

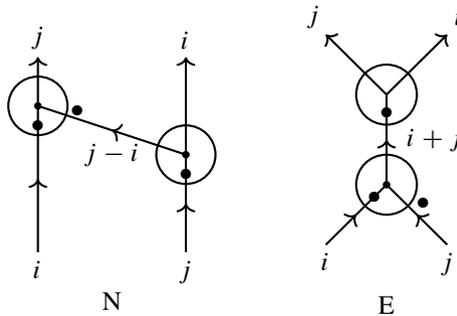
D_2 . Clearly, the factor $\mathcal{F}(\mathbb{D}, c)\mathcal{C}(D, c)$ is the same for D, D_1 and D_2 , so we are left to compare $(t^{-1/2} - t^{1/2})^{1-|V|}\langle D, c \rangle$ of both sides.

In the following calculations, we again adopt the notation $\{i\} := t^{i/2} - t^{-i/2} = [i](t^{1/2} - t^{-1/2})$. Note that a state s of the left-hand diagram D assigns one of the four corners N, E, S, W at the crossings:



The corresponding local contributions to the state sum $\langle D, c \rangle$ are $-t^{-i}, 1, t^{-i}$ and 1 , respectively.

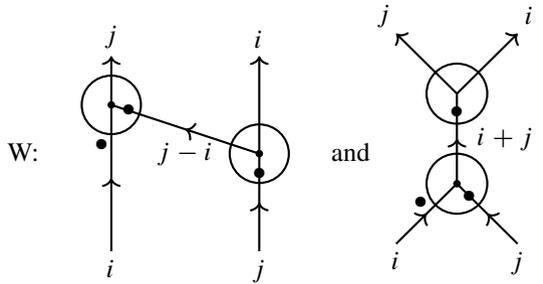
For the right-hand side, we can similarly divide all states into N-, E-, W- or S-states according to the local assignments in D_1 and D_2 . For example, there is no E-state for D_1 or N-state for D_2 ; the N-states for D_1 and the E-states for D_2 have the local assignments as depicted below,



and their contributions to the state sum are $\{i\}\{j\}t^{(j-i)/2}$ and $\{i\}\{i+j\}t^{j/2}$, respectively. Note that the number of vertices in D_1 and D_2 is 2 greater than that of D . After multiplying the coefficients in the first relation of (iv), we see that the local contributions of all N-states and E-states to both sides coincide.

$$\begin{aligned} \text{N-states:} & \quad \frac{-t^{-\frac{i+j}{2}}}{[i][j]} \cdot (t^{-1/2} - t^{1/2})^{-2} \{i\}\{j\}t^{(j-i)/2} = -t^{-i} \\ \text{E-states:} & \quad \frac{t^{-\frac{j}{2}}}{[i][i+j]} \cdot (t^{-1/2} - t^{1/2})^{-2} \{i\}\{i+j\}t^{j/2} = 1. \end{aligned}$$

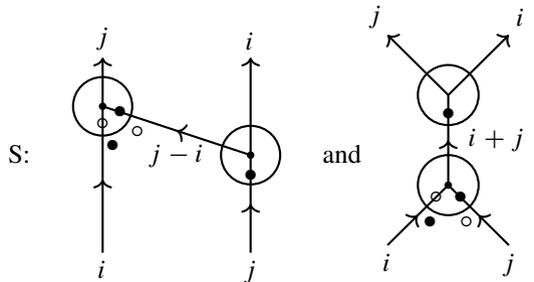
Similarly, for the local assignment (W), there are the corresponding W-states for D_1 and D_2 :



Their contributions to the state sums are $\{j - i\}\{j\}t^{-i/2}$ and $\{i + j\}\{j\}t^{-i/2}$, respectively. We can check that

$$\begin{aligned} \text{W-states: } & \frac{-t^{-\frac{i+j}{2}}}{[i][j]} \cdot (t^{-1/2} - t^{1/2})^{-2} \{j - i\}\{j\}t^{-i/2} \\ & + \frac{t^{-\frac{j}{2}}}{[i][i + j]} \cdot (t^{-1/2} - t^{1/2})^{-2} \cdot \{i + j\}\{j\}t^{-i/2} = t^{-i}. \end{aligned}$$

For the local assignment (S), there are the corresponding S-states for D_1 and D_2 :



Their contributions to the state sums sum up to $\{j\}^2$ and $\{i + j\}^2$, respectively. We can check that

$$\begin{aligned} \text{S-states: } & \frac{-t^{-\frac{i+j}{2}}}{[i][j]} \cdot (t^{-1/2} - t^{1/2})^{-2} \{j\}^2 \\ & + \frac{t^{-\frac{j}{2}}}{[i][i + j]} \cdot (t^{-1/2} - t^{1/2})^{-2} \cdot \{i + j\}^2 = 1. \end{aligned}$$

Thus, we have exhausted all states of both sides and verified the first relation in (iv).

The rest of the relations can be shown in a similar way. We remark that in the proof of relations (v) and (vii), we need to consider the change of curliness $\mathcal{C}(D, c)$ by applying Lemma 3.6. We omit the details here. \square

4.2 Graphical definition of the Alexander polynomial of a link

The MOY-type relations in the previous section provides us a graphical definition of the Alexander polynomial of a link.

Definition 4.2 Suppose D is a diagram of a link L . Let $w(D)$ be the writhe of D , i.e., the number of positive crossings  minus the number of negative crossings . Define

$$\Delta(D) := t^{w(D)/2}(\mathbb{D}), \tag{6}$$

where \mathbb{D} is the diagram D equipped with a blackboard framing and a constant coloring $c(e) = 1$ for each edge e (which corresponds to the link component in this case).

The next lemma implies that $\Delta(D)$ is an invariant of the link L .

Lemma 4.3 $\Delta(D)$ is invariant under Reidemeister moves (I), (II) and (III).

Proof For move (I), we apply the relations (ii) (iv) and (v) and obtain

$$\begin{aligned} \left(\text{Diagram of a crossing with a loop} \right) &= -t^{-1} \left(\text{Diagram of a loop with a crossing} \right) + \frac{t^{-1/2}}{[1][2]} \left(\text{Diagram of a crossing with two crossings} \right) \\ &= \frac{t^{-1/2}[1][2]}{[1][2]} \left(\text{Diagram of a crossing with two crossings} \right) = t^{-1/2} \left(\text{Diagram of a crossing with two crossings} \right). \end{aligned}$$

The remaining cases of move (I) can be proved in a similar way.

The proof of the invariance under Reidemeister moves (II) and (III) follows the same argument of [11, Theorem 3.1]. We leave the detail to the reader. \square

Note that Theorem 4.1 (iv) implies:

$$\begin{aligned} \left(\text{Diagram of a crossing with two crossings} \right) &= -t^{-1} \left(\text{Diagram of two crossings} \right) + \frac{t^{-1/2}}{[1][2]} \cdot \left(\text{Diagram of a crossing with two crossings} \right), \quad \text{and} \\ \left(\text{Diagram of a crossing with two crossings} \right) &= -t \left(\text{Diagram of two crossings} \right) + \frac{t^{1/2}}{[1][2]} \cdot \left(\text{Diagram of a crossing with two crossings} \right). \end{aligned}$$

Thus, the link invariant $\Delta(D)$ satisfies the skein relation:

$$\Delta \left(\begin{array}{c} \nearrow \quad \nwarrow \\ \nwarrow \quad \nearrow \end{array} \right) - \Delta \left(\begin{array}{c} \nwarrow \quad \nearrow \\ \nearrow \quad \nwarrow \end{array} \right) = (t^{1/2} - t^{-1/2}) \Delta \left(\begin{array}{c} \uparrow \quad \uparrow \\ \uparrow \quad \uparrow \end{array} \right).$$

In addition, Theorem 4.1 (i) implies that $\Delta(D) = 1$ when D is a trivial knot diagram. Altogether, it follows that $\Delta(D)$ is the same as the classical one-variable Alexander polynomial of a link.

4.3 MOY-type relations characterize $\Delta_{(\mathbb{G},c)}$

Grant showed in [4] that the classical MOY relations in [11] uniquely determine the MOY polynomials for colored trivalent graphs. Our next theorem gives a parallel result for the Alexander polynomial $\Delta_{(\mathbb{G},c)}(t)$.

Theorem 4.4 *The relations (i)-(x) determine $\Delta_{(\mathbb{G},c)}(t)$ for a framed trivalent graph \mathbb{G} with a positive coloring c .*

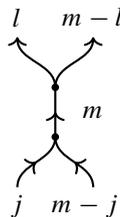
We first assert that $\Delta_{(\mathbb{G},c)}(t)$ is computable when we restrict to trivalent plane graphs colored with $\{1, 2\}$. This is analogous to a result of Grant [4, Proposition 5.2] and can be proved verbatim.

Lemma 4.5 *The relations (i)-(x) determine $\Delta_{(\mathbb{G},c)}(t)$ when (\mathbb{G}, c) is a trivalent framed plane graph colored with $\{1, 2\}$.*

Proof of Theorem 4.4 Given a framed trivalent graph diagram \mathbb{D} , we first apply (iii) and (iv) to eliminate all framing symbols and resolve all crossings to get plane graphs. Therefore, it suffices to show that $\Delta_{(\mathbb{D},c)}(t)$ is computable when \mathbb{D} is a diagram of a plane graph.

The subsequent argument is adapted from [4, Theorem 5.1]. Our goal is to use a sequence of the MOY-type relations to replace all edges colored with m by edges with colors smaller than m , where $m > 2$ is the maximal color of the graph (following the idea of Wu [14]). If this can be achieved, then for any colored diagram \mathbb{D} , there is a diagram \mathbb{D}' that is colored only with $\{1, 2\}$ such that $(\mathbb{D}) = p(t) \cdot (\mathbb{D}')$ for some function $p(t)$. It then follows from Lemma 4.5 that (\mathbb{D}) is computable.

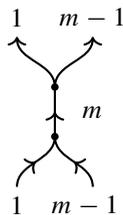
To this end, suppose that \mathbb{D} contains an edge colored with $m > 2$. Since all the edges have positive colorings, it locally looks like



with $0 < j, l < m$. If $j \geq 2$ or $l \geq 2$ we apply (vi) and (viii) to obtain the following identity:

$$\begin{aligned}
 & \left(\begin{array}{c} l \quad m-l \\ \diagdown \quad \diagup \\ \bullet \\ \diagup \quad \diagdown \\ j \quad m-j \end{array} \right) = \frac{1}{[j]^2 [l]^2} \left(\begin{array}{c} l \quad m-l \\ \diagdown \quad \diagup \\ \bullet \\ \diagup \quad \diagdown \\ j \quad m-j \end{array} \right) \\
 & \qquad \qquad \qquad = \frac{1}{[j][l][m-1]^2} \left(\begin{array}{c} l \quad m-l \\ \diagdown \quad \diagup \\ \bullet \\ \diagup \quad \diagdown \\ j \quad m-j \end{array} \right). \tag{7}
 \end{aligned}$$

Now we may assume that locally the diagram that contains the edge colored with m looks like



We then apply (ix) with $i = k = l = 1$ and $j = m - 1$ and (x) to get

$$\left(\begin{array}{c} 1 \quad m-1 \\ \swarrow \quad \searrow \\ \bullet \\ | \\ \bullet \\ \swarrow \quad \searrow \\ 1 \quad m-1 \end{array} \right) = \frac{[m]}{[1][2][m-1]} \cdot \left(\begin{array}{c} 1 \quad m-1 \\ \swarrow \quad \searrow \\ \bullet \\ \xrightarrow{1} \\ \bullet \\ \swarrow \quad \searrow \\ 1 \quad m-1 \end{array} \right) - [m-2][m] \cdot \left(\begin{array}{c} 1 \quad m-1 \\ \uparrow \quad \uparrow \\ \bullet \\ \bullet \\ \uparrow \quad \uparrow \\ 1 \quad m-1 \end{array} \right).$$

Observe that the right-hand side diagrams of the equation has no local colors larger than $m - 1$ now. This is exactly what we aimed for and thus concludes the proof. \square

5 Properties and applications

In this section, we apply the state sum formula and the MOY-type relations established in the previous sections to prove some basic properties of the Alexander polynomials $\Delta_{(G,c)}(t)$ and $\Delta_{(\mathbb{G},c)}(t)$. As an application, we show that the Alexander polynomial can be used to detect the non-planarity and chirality of a spatial graph.

5.1 Symmetries

Suppose D is a diagram of a graph G . Recall that the *mirror image* of D , denoted D^* , is obtained from D by changing all positive (resp. negative) crossings of D into negative (resp. positive) ones. The *orientation reversal* of D , denoted $-D$, is obtained from reversing all the edge orientations of D . We define G^* and $-G$ as the graphs represented by D^* and $-D$, respectively. It is standard that G^* and $-G$ only depends on G and not on D . When we have a framed graph \mathbb{G} , we can define \mathbb{G}^* and $-\mathbb{G}$ in a similar way by also changing the framing symbols $\neg\curvearrowright$ and $\neg\curvearrowleft$ in the corresponding diagrams accordingly.

Proposition 5.1 *The normalized Alexander polynomial $\Delta_{(\mathbb{G},c)}(t)$ of a framed trivalent graph with positive coloring satisfies the following symmetric properties.*

- (i) $\Delta_{(\mathbb{G},c)}(t) = \Delta_{(\mathbb{G}^*,c)}(t^{-1})$.
- (ii) $\Delta_{(\mathbb{G},c)}(t) = \Delta_{(-\mathbb{G},c)}(t)$.

In particular, a plane graph \mathbb{G} with the blackboard framing satisfies the symmetric relation

$$\Delta_{(\mathbb{G},c)}(t) = \Delta_{(\mathbb{G},c)}(t^{-1}).$$

Proof (i) Given a framed diagram \mathbb{D} of \mathbb{G} , define

$$\tilde{\Delta}_{(\mathbb{D},c)}(t) := \frac{\Delta_{(\mathbb{D},c)}(t) + \Delta_{(\mathbb{D}^*,c)}(t^{-1})}{2}.$$

One can check that $\tilde{\Delta}$ satisfy all the MOY-type relations in Theorem 4.1. On the other hand, we know from Theorem 4.4 that those relations completely characterize the Alexander polynomial. In other words, $\tilde{\Delta}_{(\mathbb{D},c)}(t) = \Delta_{(\mathbb{D},c)}(t)$. Hence, $\Delta_{(\mathbb{D},c)}(t) = \Delta_{(\mathbb{D}^*,c)}(t^{-1})$.

(ii) We can argue in a similar way. This time let

$$\tilde{\Delta}_{(\mathbb{D},c)}(t) := \frac{\Delta_{(\mathbb{D},c)}(t) + \Delta_{(-\mathbb{D},c)}(t)}{2},$$

and we are going to check that $\tilde{\Delta}$ satisfy all the MOY-type relations in Theorem 4.1. Note that all relations except (vii) and (ix) are obviously true with the oppositely oriented edges, so it suffices for us to show that (vii) and (ix) also hold after the reverse of orientations. More precisely, we need to verify the relations

$$\begin{aligned} \left(\begin{array}{c} j \quad i+j \quad i \\ \curvearrowright \quad \curvearrowright \quad \curvearrowright \\ i \quad \quad \quad j \\ \curvearrowleft \quad \curvearrowleft \quad \curvearrowleft \\ j \quad i+j \quad i \end{array} \right) &= \frac{[i+j]^3}{[i]} \cdot \left(\begin{array}{c} j \quad \quad \quad i \\ \curvearrowright \quad \quad \quad \curvearrowright \\ \quad \quad \quad i-j \\ \curvearrowleft \quad \quad \quad \curvearrowleft \\ j \quad \quad \quad i \end{array} \right) \\ &+ [j]^2 [i+j]^2 \cdot \left(\begin{array}{c} \downarrow \quad \quad \quad \uparrow \\ j \quad \quad \quad i \end{array} \right), \end{aligned}$$

and

$$\begin{aligned} \left(\begin{array}{c} i+k-l \quad j+l-k \\ \downarrow \quad \quad \quad \downarrow \\ \quad \quad \quad l \\ \downarrow \quad \quad \quad \downarrow \\ i+k \quad j-k \\ \downarrow \quad \quad \quad \downarrow \\ i \quad \quad \quad j \\ \quad \quad \quad k \end{array} \right) &= \frac{[j][l][i+k]}{[i+j]} \cdot \left(\begin{array}{c} i+k-l \quad j+l-k \\ \curvearrowright \quad \quad \quad \curvearrowright \\ \quad \quad \quad i+j \\ \downarrow \quad \quad \quad \downarrow \\ i \quad \quad \quad j \end{array} \right) \\ &+ [i+k][j-k] \cdot \left(\begin{array}{c} i+k-l \quad j+l-k \\ \downarrow \quad \quad \quad \downarrow \\ \quad \quad \quad k-l \\ \downarrow \quad \quad \quad \downarrow \\ i \quad \quad \quad j \end{array} \right). \end{aligned}$$

As both identities can be proved in the same way as the relations in Theorem 4.1, we omit the details here.

On the other hand, we know from Theorem 4.4 that those relations completely characterize the Alexander polynomial. In other words, $\tilde{\Delta}_{(\mathbb{D},c)}(t) = \Delta_{(\mathbb{D},c)}(t)$. Hence, $\Delta_{(\mathbb{D},c)}(t) = \Delta_{(-\mathbb{D},c)}(t)$.

5.2 Integrality

Recall from Definition 3.2 that the Alexander polynomial $\Delta_{(G,c)}(t)$ equals $\frac{\langle D,c \rangle}{(t^{-1/2} - t^{1/2})^{|V|-1}}$ for any diagram D of G . A priori, $\Delta_{(G,c)}(t)$ is only a rational function of the variable $t^{\pm \frac{1}{2}}$. Our next proposition shows that $\Delta_{(G,c)}(t)$ is a genuine polynomial of $t^{\pm \frac{1}{2}}$ when G has at least one vertex.

Proposition 5.2 *Let G be an MOY graph with at least one vertex. Then the Alexander polynomial $\Delta_{(G,c)}(t) \in \mathbb{Z}[t^{\pm \frac{1}{2}}]$.*

Proof Recall that $\langle D, c \rangle = |\delta|^{-1} \sum_{s \in S(D,\delta)} M(s) \cdot A(s)$, where $A(s) = \prod_{p=1}^N A_{C_p}^{s(C_p)}$ is defined by the local contribution as exhibited in Fig. 6 (bottom). Our key observation is that the contribution $A_{C_p}^{s(C_p)}$ for a state s at the crossing C_p of type  assigning $s(C_p)$ the north corner inside the circular region is equal to $t^{-i/2} - t^{i/2}$ for some i , which is a factor divisible by $t^{-1/2} - t^{1/2}$.

To compute $\langle D, c \rangle$ for the graph with $|V| \geq 1$ vertices, we place the base point δ on an edge e_1 of color i_1 just before it goes into a vertex v_1 . Then, all states in this decorated diagram must assign the north corner inside the circular region corresponding to v_1 , which contributes a factor $A_{C_1}^\Delta = t^{-i_1/2} - t^{i_1/2}$. As

$$|\delta| = t^{n-i_1} - t^n = t^{n-i_1/2} \cdot (t^{-i_1/2} - t^{i_1/2})$$

for some n , the factor $|\delta|$ is cancelled by $A_{C_1}^\Delta$. In addition, note that $A(s)$ of each state s consists of a total of $|V| - 1$ additional factors of the form $t^{-i/2} - t^{i/2}$, each of which corresponds to the contribution of $A_{C_p}^{s(C_p)}$ coming from the remaining $|V| - 1$ circular regions. Consequently, $\frac{A(s)}{|\delta|(t^{-1/2} - t^{1/2})^{|V|-1}} \in \mathbb{Z}[t^{\pm \frac{1}{2}}]$ for all states s , and so is the Alexander polynomial $\Delta_{(G,c)}(t)$. \square

5.3 Positivity

One of the fundamental problems in classical knot theory is to determine whether a given diagram is a projection of the unknot. The Alexander polynomial of knots gives a simple primary test for the triviality of knots. Analogously, an interesting problem in spatial graph theory is to determine whether a given diagram is isotopic to a plane graph diagram. We give in this section a necessary condition in terms of the Alexander polynomial of graphs.

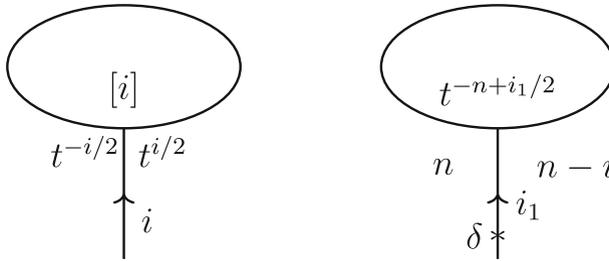


Fig. 13 The local contributions $P_{C_p}^\Delta$ for crossings that correspond to generic edges (left) and the marked edge (right), respectively

When D is a plane graph diagram, we first reformulate the state sum formula of $\Delta_{(D,c)}(t)$. To this end, we have only the crossings of type \updownarrow and we define a new local contribution $P_{C_p}^\Delta$ as in Fig. 13.

For each state s , let $P(s) = \prod_{p=1}^N P_{C_p}^{s(C_p)}$. Note that for the crossing C_1 corresponding to the edge e_1 where δ lies,

$$P_{C_1}^{s(C_1)} = t^{-n+i_1/2} = \frac{t^{-i_1/2} - t^{i_1/2}}{t^{n-i_1} - t^n} = \frac{M_{C_1}^{s(C_1)} \cdot A_{C_1}^{s(C_1)}}{|\delta|};$$

and for the north corner $s(C_p)$ of a generic crossing C_p ,

$$P_{C_p}^{s(C_p)} = [i] = \frac{t^{-i/2} - t^{i/2}}{t^{-1/2} - t^{1/2}} = \frac{1}{t^{-1/2} - t^{1/2}} \cdot M_{C_p}^{s(C_p)} \cdot A_{C_p}^{s(C_p)}.$$

The proof of Proposition 5.2 shows that

$$\sum_{s \in S(D,\delta)} P(s) = |\delta|^{-1} \frac{\sum_{s \in S(D,\delta)} M(s) \cdot A(s)}{(t^{-1/2} - t^{1/2})^{|V|-1}}$$

Hence,

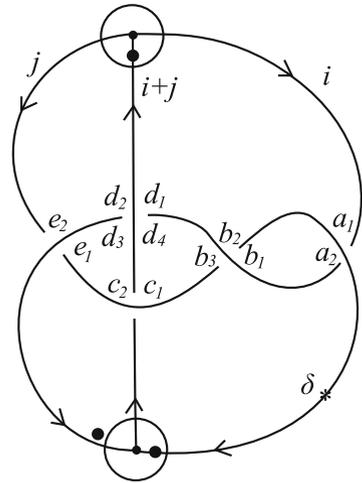
$$\Delta_{(D,c)}(t) = \sum_{s \in S(D,\delta)} P(s). \tag{8}$$

Using this new version of the state sum formula for the Alexander polynomial, it is straightforward to see $\Delta_{(G,c)}(t) \in \mathbb{Z}[t^{\pm 1/2}]$, but we can actually prove more. Note that

$$[i] = \frac{t^{i/2} - t^{-i/2}}{t^{1/2} - t^{-1/2}} = t^{\frac{i-1}{2}} + \dots + t^{\frac{1-i}{2}}$$

is a polynomial of positive coefficient when $i > 0$; so are all the other local contribution $P_{C_p}^\Delta$. Thus, every term $P(s)$ is a polynomial of positive coefficient. Summing up, we obtain the following positivity condition on the Alexander polynomial of a plane graph.

Fig. 14 The 5_1 in Litherland's table of θ -curve diagrams. Observe that any two of three curves of the graph form an unknot. With the given position of δ , the local assignment of all states at the circular regions is unique and is marked out by big black dots



Theorem 5.3 Suppose G is isotopic to a plane MOY graph with at least one vertex. Then $\Delta_{(G,c)}(t) \in \mathbb{Z}_{\geq 0}[t^{\pm \frac{1}{2}}]$.

Remark 5.4 Note that the ambiguity of a power of t in the definition of $\Delta_{(G,c)}(t)$ does not matter, since the sign of the coefficients is unchanged.

Example 5.5 We conclude this section with a detailed calculation of the graph G exhibit in Fig. 14, whose undirected version is denoted 5_1 in Litherland's table of θ -curve diagrams [10, Fig. 1]. In this example, we show the non-planarity and chirality of this graph.

In the following calculation, we assign the balanced colors i, j and $i + j$ and place the base point δ as in Fig. 14. We want to compute $\Delta_{(G,c)}(t)$ using the state sum formula, so we begin by enumerating all Kauffman states in the given decorate diagram. Note that the assignment at the 3 crossings of type \downarrow is unique; thus, in order to specify the states, it is enough to describe the local assignment at the 5 crossings of types \nearrow and \nwarrow . By carefully exhausting all possibilities, we find 7 states in total. We then compute the value of $\frac{M(s) \cdot A(s)}{|\delta|(t^{-1/2} - t^{1/2})^{|V|-1}}$ for each state s and summarize the result in Table 2.

Therefore,

$$\begin{aligned} \Delta_{(G,c)}(t) &= \sum_{s \in S(D,\delta)} \frac{M(s) \cdot A(s)}{|\delta|(t^{-1/2} - t^{1/2})^{|V|-1}} \\ &= \left(t^{\frac{-3i-3j}{2}} - t^{\frac{-3i-j}{2}} - t^{\frac{-i-3j}{2}} + t^{\frac{-i-j}{2}} + t^{\frac{i-j}{2}} + t^{\frac{j-i}{2}} - t^{\frac{i+j}{2}} \right) \cdot [i + j] \end{aligned}$$

Suppose $i, j > 0$. Since $[i + j] = t^{\frac{i+j-1}{2}} + \dots + t^{\frac{1-i-j}{2}}$, the term of the highest degree in $\Delta_{(G,c)}(t)$ is $-t^{-\frac{1}{2}+i+j}$, which has a negative coefficient -1 . By Theorem 5.3, the given graph 5_1 in Fig. 14 is not isotopic to any plane graph.

Table 2 All states and their values of $\frac{M(s) \cdot A(s)}{|\delta|(t^{-1/2} - t^{1/2})^{|V|-1}}$

State s	Value of $\frac{M(s) \cdot A(s)}{ \delta (t^{-1/2} - t^{1/2})^{ V -1}}$
$a_1 b_1 c_1 d_2 e_1$	$t^{\frac{-3i-3j}{2}} \cdot [i + j]$
$a_1 b_1 c_1 d_3 e_2$	$-t^{\frac{-i-3j}{2}} \cdot [i + j]$
$a_1 b_1 c_2 d_4 e_2$	$t^{\frac{i-j}{2}} \cdot [i + j]$
$a_2 b_2 c_1 d_2 e_1$	$-t^{\frac{-3i-j}{2}} \cdot [i + j]$
$a_2 b_2 c_1 d_3 e_2$	$t^{\frac{-i-j}{2}} \cdot [i + j]$
$a_2 b_2 c_2 d_4 e_2$	$-t^{\frac{+i+j}{2}} \cdot [i + j]$
$a_2 b_3 c_2 d_1 e_2$	$t^{\frac{j-i}{2}} \cdot [i + j]$

Furthermore, if we set the color $i = j = 1$, the Alexander polynomial

$$\Delta_{(G,c)}(t) = (t^{-3} - 2t^{-2} + t^{-1} + 2 - t) \cdot (t^{\frac{1}{2}} + t^{-\frac{1}{2}}).$$

As $\Delta_{(G,c)}(t) \neq \Delta_{(G,c)}(t^{-1}) \cdot t^k$ for any k , Proposition 5.1 in fact implies the stronger result that S_1 is *chiral*, i.e., the graph is not isotopic to its mirror.

5.4 Non-vanishing properties

In this section, our main goal is to establish the non-vanishing property for the Alexander polynomial of a connected trivalent plane graph with a positive coloring. One should compare it with Kronheimer-Mrowka [8, Theorem 1.1], although it is at present unclear to the authors about the connection of the Alexander polynomial to Tait colorings.

Theorem 5.6 *Suppose G is a connected trivalent plane graph with a positive coloring c . Then $\Delta_{(G,c)}(t) \neq 0$.*

Recall that a *bridge* is an edge of a graph whose removal increases its number of connected components. Note that the color of a bridge has to be 0 for a balanced coloring. Hence, the graph G that has a positive coloring in Theorem 5.6 must not have a bridge. In contrast, we claim that $\Delta_{(G,c)}(t)$ vanishes when G is either not connected or contains a bridge.

Proposition 5.7 *Suppose G is a plane graph that is either not connected or contains a bridge. Then $\Delta_{(G,c)}(t) = 0$.*

Proof Suppose D is a diagram of G . Remember that the state sum $\langle D, c \rangle$ is defined to be 0 when D is not connected. On the other hand, if D has a bridge, then the color of the edge must be 0. We can apply the relation (x) in Theorem 4.1 to obtain a disconnected diagram whose Alexander polynomial is 0. □

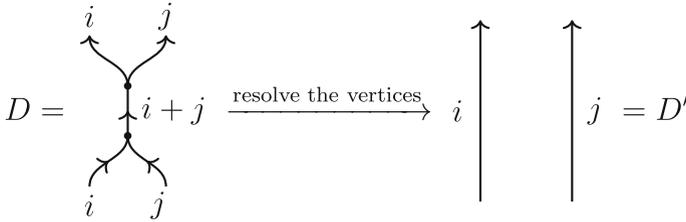


Fig. 15 Resolve the vertices of a local configuration

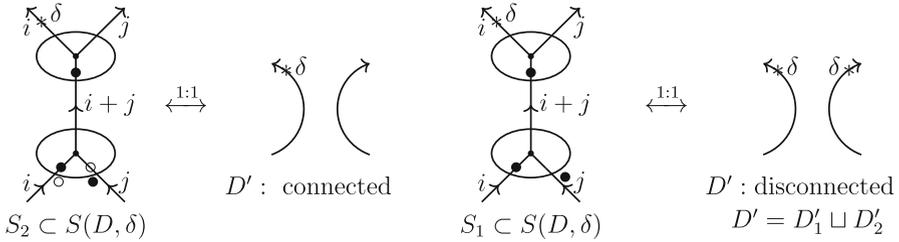


Fig. 16 Splitting the set of states $S(D, \delta)$ to S_1 and S_2 . Note that there are natural correspondences $S_1 \xleftrightarrow{1:1} S(D'_1, \delta) \times S(D'_2, \delta)$ when $D' = D'_1 \sqcup D'_2$, and $S_2 \xleftrightarrow{1:1} S(D', \delta)$ when D' is connected

The rest of the section is devoted to the proof of Theorem 5.6. When the vertex set V of G is empty, the graph G is simply a trivial knot. In that case, $\Delta_{(G,c)}(t)$ is non-zero by Theorem 4.1(i). From now on, let us assume that D is a plane graph diagram of G with at least one vertex with a positive coloring c , unless otherwise specified.

Lemma 5.8 *Suppose D is a trivalent plane graph diagram with a positive coloring c . Then $\Delta_{(D,c)}(t) \neq 0$ if and only if the set of states $S(D, \delta) \neq \emptyset$.*

Proof This follows from Equation (8) and the observation that each term $P(s)$ is a polynomial of positive coefficients. \square

Our plan is to prove $S(D, \delta) \neq \emptyset$ by induction. Note that from Lemma 5.8, the set $S(D, \delta) \neq \emptyset$ if and only if $S(D, \delta') \neq \emptyset$ for any two points δ and δ' . Thus we have the freedom of placing the base point anywhere on D . The next step is to relate the set $S(D, \delta)$ to the set $S(D', \delta)$ of a simpler graph diagram D' . To this end, we start with a graph diagram D that contains a local configuration as Fig. 15 and resolve the vertices to get a new graph diagram D' . Clearly, D' is also a trivalent plane graph with a positive coloring. In some cases, D' may be disconnected. Then decompose $D' = D'_1 \sqcup D'_2$ so that the graphs D'_1 and D'_2 are connected.

Lemma 5.9 *For D, D' as above, if D' is connected, then $S(D', \delta) \neq \emptyset$ implies that $S(D, \delta) \neq \emptyset$; if D' is disconnected, then $S(D'_i, \delta) \neq \emptyset$ for both $i = 1, 2$ implies that $S(D, \delta) \neq \emptyset$.*

Proof We split the set of states (if there are any) $S(D, \delta) = S_1 \sqcup S_2$ according to the local assignment shown in Fig. 16.

Suppose the resulted graph D' from the above resolution is connected. We claim that $|S_2| = |S(D', \delta)|$, because every state in $S(D', \delta)$ naturally extends to a state in S_2 . In particular, $S(D', \delta) \neq \emptyset$ implies that $S(D, \delta) \neq \emptyset$.

Otherwise, D' is disconnected and assume that $D' = D'_1 \sqcup D'_2$. We claim that $|S_1| = |S(D'_1, \delta)| \times |S(D'_2, \delta)|$, because a pair of states in $S(D'_i, \delta)$ can be naturally extended to a state in S_1 . In particular, $S(D'_i, \delta) \neq \emptyset$ for $i = 1, 2$ implies that $S(D, \delta) \neq \emptyset$. \square

Lemma 5.10 *Suppose D is a connected trivalent plane graph diagram with at least one vertex, and there is a balanced color c so that each edge is colored with $\{1, 2\}$. Then $S(D, \delta) \neq \emptyset$.*

Proof We prove by induction on the number of vertices $|V|$ of D . Since D is trivalent, $|V|$ is an even number. The base case is $|V| = 2$. The trivial θ -curve is the unique planar trivalent graph of 2 vertices. In that case, $|S(D, \delta)| = 1$.

Suppose that the statement is true for all $|V| \leq 2n$. We now consider a graph diagram D with $2n + 2$ vertices. Note that D must contain a local configuration of Fig. 15 with $i = j = 1$, so we resolve the vertices and obtain D' . Since D' (or D'_1 and D'_2 in the case when D' is disconnected) is a connected trivalent plane graph diagram with a strictly fewer number of vertices, we have the inductive hypothesis that $S(D', \delta) \neq \emptyset$ (resp. $S(D'_i, \delta) \neq \emptyset$ for $i = 1, 2$ when D' is disconnected). Hence, Lemma 5.9 implies that $S(D, \delta) \neq \emptyset$. \square

Proof of Theorem 5.6 We prove by induction on the maximal color m on the graph G , which is the same type of argument we used in proving Theorem 4.4. When $m = 1$, the vertex set V is empty and we have already discussed the case. When $m = 2$, the statement follows from Lemma 5.8 and Lemma 5.10. Thus it suffices to consider $m \geq 3$.

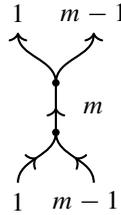
Suppose now that $\Delta_{(D,c)}(t) \neq 0$, or equivalently, $S(D, \delta) \neq \emptyset$ for all graph diagram D with maximal color $m - 1$. We want to prove the same thing for a graph diagram D with maximal color m .

First consider the case when there is a unique edge on D with the maximal color m . From Equation (7), we see that

$$\left(\begin{array}{c} l \quad m-l \\ \diagup \quad \diagdown \\ \bullet \\ \diagdown \quad \diagup \\ m \\ \bullet \\ \diagup \quad \diagdown \\ j \quad m-j \end{array} \right) = \frac{1}{[j][l][m-1]^2} \left(\begin{array}{c} l \quad m-l \\ \diagup \quad \diagdown \\ \bullet \\ \leftarrow 1 \quad \rightarrow \\ \bullet \\ \diagdown \quad \diagup \\ m-1 \\ \bullet \\ \diagdown \quad \diagup \\ m \\ \bullet \\ \diagup \quad \diagdown \\ \rightarrow 1 \quad \leftarrow \\ \bullet \\ \diagdown \quad \diagup \\ j \quad m-j \end{array} \right).$$

Since the Alexander polynomial of the two diagrams are either simultaneously zero or simultaneously non-zero, it suffices to prove the non-vanishing property of a diagram

D on the right-hand side of the equation that contains the local configuration:



Now, we can smooth the crossing and obtain D' as in Fig. 15. Since the maximal color of D' (or D'_1 and D'_2 in the case when D' is disconnected) is at most $m - 1$, we have $S(D', \delta) \neq \emptyset$ (resp. $S(D'_i, \delta) \neq \emptyset$ for $i = 1, 2$ when D' is disconnected) by induction. Hence, Lemma 5.9 implies that $S(D, \delta) \neq \emptyset$, and Lemma 5.8 implies that $\Delta_{(D,c)}(t) \neq 0$.

In general, suppose D has k edges of the maximal color m . We can apply another induction on k and repeat the above argument to reduce k . This concludes the proof. □

5.5 An intrinsic invariant

So far, we have implicitly assumed *embedded graphs* in this paper, i.e., they are the graphs that exist in a specific position. In contrast, an *abstract graph* is a graph that is considered to be independent of any particular embedding. Whereas an embedded graph G has a unique underlying abstract graph g , a given abstract graph g can be typically embedded in many different ways and gives rise to distinct embedded graphs. A property or invariant of a graph is called *intrinsic* if it only depends on the underlying abstract graph (and is independent of the embedding).

Suppose G is an MOY graph. Recall that (3) defines $\Delta_{(G,c)}(t)$ up to a power of t ; hence, $\Delta_{(G,c)}(1)$ is a well-defined invariant that is independent of the graph diagrams. Our next result shows that $\Delta_{(G,c)}(1)$ is furthermore an intrinsic invariant of graphs. This may be viewed as a generalization of the classical result that the Alexander polynomial of a knot satisfies $\Delta_K(1) = 1$.

Proposition 5.11 *The value of the Alexander polynomial evaluated at $t = 1$, $\Delta_{(G,c)}(1)$, is an integer-valued invariant of the underlying abstract MOY graph g .*

Proof Note that the local contributions $M_{C_p}^\Delta$ and $A_{C_p}^\Delta$ evaluated at $t = 1$ is invariant under crossing changes (Compare the first and second column of Fig. 6). Since any other embedded graphs G' with the same underlying abstract graph g can be obtained from G via the Reidemeister moves in Fig. 2 and crossing changes, it follows that $\Delta_{(G,c)}(1) = \Delta_{(G',c)}(1)$; in other words, it is indeed an intrinsic invariant of the abstract MOY graph g . □

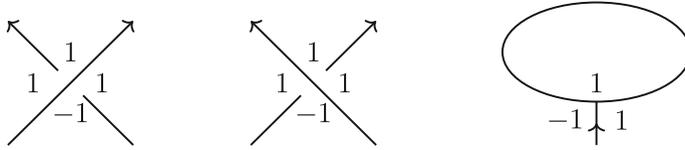


Fig. 17 The local contribution of $\text{sign}_{C_p}^\Delta$

In light of Proposition 5.11, we propose to define an *Alexander invariant* for an abstract MOY graph g by letting

$$\Delta_{(g,c)} := \Delta_{(G,c)}(1), \tag{9}$$

where G is an arbitrary MOY embedding of g .

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Appendix A: Replacing $\text{sign}(s) \cdot m(s)$ with $M(s)$

The goal of this appendix is to show Proposition 2.16. To this end, it suffices to prove the following Proposition. Let (D, δ) be a decorated MOY graph diagram. For any given total orders on $\text{Cr}(D)$ and $\text{Re}(D)$, recall that $\text{sign}(s)$ is defined to be the sign of the state $s \in S(D, \delta)$ as a permutation with respect to the given orders.

Proposition A.1 *Let (D, δ) be a decorated MOY graph diagram, we have*

$$\frac{\text{sign}(s_1)}{\text{sign}(s_2)} = \frac{M(s_1)m(s_1)}{M(s_2)m(s_2)}, \tag{10}$$

for any two states $s_1, s_2 \in S(D, \delta)$. Namely we have $\text{sign}(s) = \prod_{p=1}^N \text{sign}_{C_p}^{s(C_p)}$ up to an overall sign change, where $\text{sign}_{C_p}^{s(C_p)}$ is the local contribution defined in Fig. 17.

Proof It is easy to see that any diagram D can be transformed into a trivalent graph diagram by surgery (ΛI) or $(I \Lambda)$ and surgery (I) , which are defined in Figs. 22 and 21. Therefore the proof is a combination of Lemmas A.3, A.4 and A.5. \square

Lemma A.2 *Let D be a diagram of a singular link. Namely at each vertex, there are two edges pointing inwards and two outwards. Then the statement of Theorem A.1 holds.*

The Heegaard Floer homology and the Alexander polynomial of a singular link are studied in [12]. Using statements there we can get Lemma A.2. However to keep the combinatorial flavor of the paper, we give a proof by recalling some facts in [6].

A universe U is a connected diagram of a link in \mathbb{R}^2 where the data of which strand is over and which is under at each crossing is suppressed. Choose a base point δ on U and consider the set of states $S(U, \delta)$. The Clock Theorem in [6] says that any two states in $S(U, \delta)$ are connected by a sequence of clockwise or counterclockwise transpositions, which are defined in Fig. 18.

Two states that are connected by a transposition satisfy (10). By the Clock Theorem, any two states in $S(U, \delta)$ satisfy (10). Therefore the sign of a state for a universe is given by the local contribution in Fig. 17.

Proof of Lemma A.2 Since D is a connected diagram of a singular link, we call the underlying universe U_D . The difference of $\text{Cr}(D)$ and $\text{Cr}(U_D)$ occurs around the singular crossings, and the same is true for $\text{Re}(D)$ and $\text{Re}(U_D)$.

Suppose C_i is a crossing of U_D which corresponds to a singular crossing of D . There are two crossings in $\text{Cr}(D)$ and one circle region in $\text{Re}(D)$ around C_i , which we call $C_i, C_{i.5}$ and $R_{i.5}$, as in Fig. 19. We choose orders on $\text{Cr}(D)$ and $\text{Re}(D)$ so that

$$C_i < C_{i.5} < C_{i+1} \text{ and } R_i < R_{i.5} < R_{i+1}.$$

Then by forgetting $C_{i.5}$ and $R_{i.5}$ we get orders on $\text{Cr}(U_D)$ and $\text{Re}(U_D)$.

We see that any state s in $S(D, \delta)$ corresponds to a state \tilde{s} in $S(U_D, \delta)$ by ignoring the circle regions:

$$\tilde{s}(C_i) = \begin{cases} \text{souther corner,} & s(C_i) = \text{east corner} \vee s(C_{i.5}) = \text{west corner} \\ \text{west corner,} & s(C_i) = \text{west corner} \\ \text{east corner,} & s(C_i) = \text{east corner} \end{cases}$$

Considering the orders above, we have

$$\text{sign}(s) = \begin{cases} \text{sign}(\tilde{s}), & \text{if } s(C_{i.5}) = R_{i.5} \\ -\text{sign}(\tilde{s}), & \text{if } s(C_i) = R_{i.5}. \end{cases}$$

Up to an overall sign change, it is easy to check that the signs defined from the local contributions in Fig. 19 satisfies the relation above. □

Lemma A.3 *When D is a diagram of an oriented trivalent graph without sinks or sources, the statement of Theorem A.1 holds.*

To prove Lemma A.3, we separate the vertices of a trivalent graph into two groups. We call a vertex with indegree two and outdegree one an *even vertex*, and a vertex

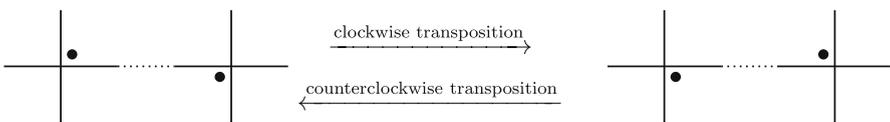


Fig. 18 Transpositions in the Clock Theorem

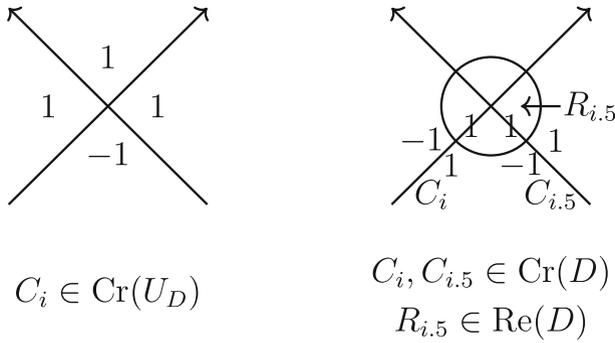


Fig. 19 The local contribution of sign for a universe U_D (left) and a graph D (right)

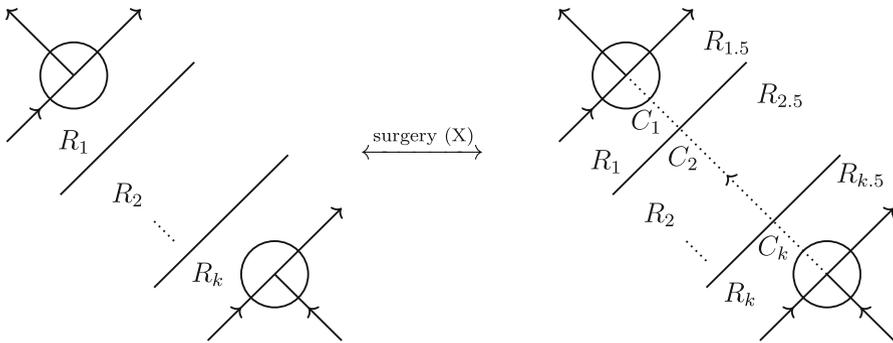


Fig. 20 The surgery (X) transforms a trivalent graph into a singular link

with indegree one and outdegree two an *odd vertex*. It is easy to see that the number of even vertices equals that of odd vertices. We use an oriented simple arc to connect an even vertex to an odd vertex, and call it surgery (X) , which transforms two trivalent vertices into singular crossings.

Proof of Lemma A.3 Let D_1 and D_2 be the graph diagrams before and after applying a surgery (X) . We show that if the statement of Theorem A.1 holds for D_2 , it also holds for D_1 .

As shown in Fig. 20, the interior of the newly added arc intersects D_1 at several points. We assume that the regions that the arc goes across, which we call R_1, R_2, \dots, R_k , are distinct regions, otherwise we can replace the arc by an arc with less intersection points with D_1 .

The arc separates each region R_q for $1 \leq q \leq k$ into two regions since D_1 is a connected diagram. We label the regions in D_2 around the arc by $R_1, R_{1.5}, R_2, R_{2.5}, \dots, R_k, R_{k.5}$ as shown in Fig. 20, and label the newly created crossings by C_1, C_2, \dots, C_k . Note that $\text{Cr}(D_2) = \text{Cr}(D_1) \cup \{C_1, C_2, \dots, C_k\}$. Consider an order on $\text{Cr}(D_1)$ and extend it to an order on $\text{Cr}(D_2)$ by requiring that

$$C_1 < C_2 < \dots < C_k < \text{any other crossing in } \text{Cr}(D_1).$$

Consider an order on $\text{Re}(D_2)$ so that

$$R_1 < R_{1.5} < R_2 < R_{2.5} < \dots < R_k < R_{k.5}.$$

Then we get an order on $\text{Re}(D_1)$ by forgetting $R_{i.5}$'s.

Surgery (X) naturally induces an injective map $\phi : S(D_1, \delta) \rightarrow S(D_2, \delta)$ described as below. Given $s \in S(D_1, \delta)$, $\phi(s)$ sends the crossing C_p to the region R_p (resp. $R_{p.5}$) if s does not occupy any corner in R_p (resp. $R_{p.5}$), for $1 \leq p \leq k$.

By considering the orders above, we can check that

$$\frac{\text{sign}(s_1)}{\text{sign}(s_2)} = \frac{\text{sign}(\phi(s_1)) \prod_{p=1}^k \text{sign}_{C_p}^{\phi(s_1)(C_p)}}{\text{sign}(\phi(s_2)) \prod_{p=1}^k \text{sign}_{C_p}^{\phi(s_2)(C_p)}} = \frac{\prod_{C_p > C_k} \text{sign}_{C_p}^{\phi(s_1)(C_p)}}{\prod_{C_p > C_k} \text{sign}_{C_p}^{\phi(s_2)(C_p)}}$$

for any two states $s_1, s_2 \in S(D_1, \delta)$. The first equality comes from a direct calculation of the sign using the orders defined above, and the second equality follows from the assumption that the statement of Theorem A.1 holds for D_2 . Note that the crossings of D_2 that are greater than C_k are exactly the crossings of D_1 . Therefore

$$\frac{\text{sign}(s_1)}{\text{sign}(s_2)} = \frac{\prod_{C_p \in \text{Cr}(D_1)} \text{sign}_{C_p}^{s_1(C_p)}}{\prod_{C_p \in \text{Cr}(D_1)} \text{sign}_{C_p}^{s_2(C_p)}}.$$

This completes the proof of the lemma. □

Lemma A.4 *If the statement of Theorem A.1 holds for the diagram after a surgery (I), then it also holds for the one before the surgery (I) (see Fig. 21).*

Proof Suppose the diagrams before and after a surgery (I) are D_1 and D_2 . It is easy to see that the surgery (I) induces a one-one map $\phi : s(D_1, \delta) \rightarrow s(D_2, \delta)$ since C_1 must be mapped to R_1 . Therefore

$$\frac{\text{sign}(s_1)}{\text{sign}(s_2)} = \frac{\text{sign}(\phi(s_1))}{\text{sign}(\phi(s_2))},$$

the right-hand side of which, by assumption, is defined by the local contribution of sign. This completes the proof. □

Lemma A.5 *If the statement of Theorem A.1 holds for the diagram after a surgery (ΔI) or ($I \Delta$), then it also holds for the one before the surgery (ΔI) or ($I \Delta$) (see Fig. 22).*

Proof We prove the lemma for surgery (ΔI), and the case of surgery ($I \Delta$) can be proved similarly. Suppose the diagrams before and after surgery (ΔI) are D_1 and D_2 . Consider an order on $\text{Cr}(D_2)$, which induces an ordering in $\text{Cr}(D_1)$ by forgetting C_1 . Since $\text{Re}(D_2) = \text{Re}(D_1) \cup \{R_1\}$, an order on $\text{Re}(D_2)$ also induces an order on $\text{Re}(D_1)$ by forgetting R_1 .

Fig. 21 Surgery (I)

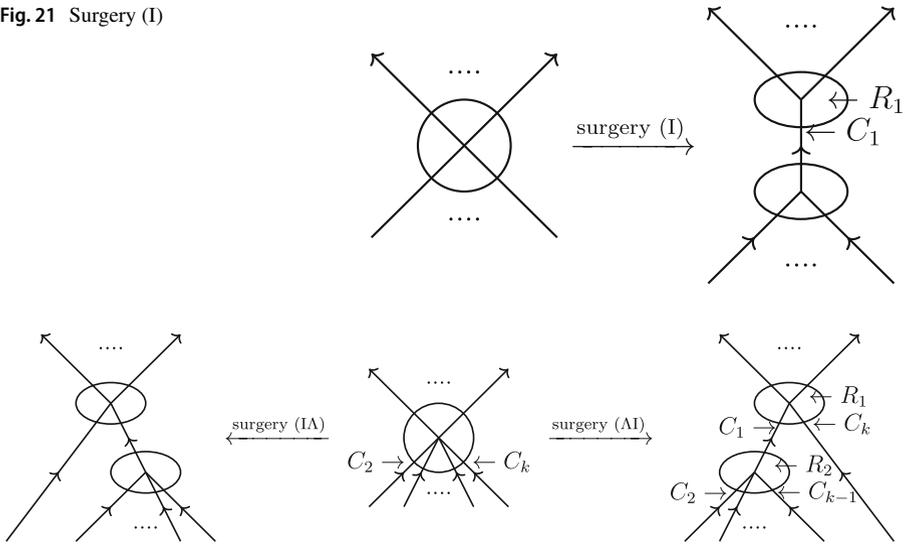


Fig. 22 Surgery (IA) and (AI)

The surgery (AI) induces an injective map $\phi : s(D_1, \delta) \rightarrow s(D_2, \delta)$ as below. Given $s \in s(D_1, \delta)$, if $s(C_k) \neq R_2$, let $\phi(s)(C_1) = R_1$ and $\phi(s)$ maps the other crossings the same way as s . If $s(C_k) = R_2$ and $s(C_{k-1})$ is the east corner of C_{k-1} , let $\phi(s)(C_k) = R_1$, $\phi(s)(C_{k-1}) = R_2$ and $\phi(s)(C_1)$ be the east corner. If $s(C_k) = R_2$ and $s(C_{k-1})$ is the west corner, then $s(C_2)$ must be its west corner. In this case, let $\phi(s)(C_k) = R_1$, $\phi(s)(C_2) = R_2$ and $\phi(s)(C_1)$ the west corner.

In each case, we can check that $\text{sign}(s) = \text{sign}(\phi(s))$ for any $s \in s(D_1, \delta)$. The right-hand side of the equality, by assumption, is defined by the local contribution of sign . Since ϕ is an injective map, the local definition of sign works as well for D_1 . \square

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