Lecture 8
Example: Consider:
$$a \frac{d^2u}{dx^2} + b \frac{du}{dx} = f$$
 for $x \in [0, 2\pi]$.
This time, we approximate $\frac{d^2u}{dx^2}$ by =
(*) $\frac{d^2u}{dx^2}(x_j) \approx \frac{u_{j-2} - 2u_{j} + u_{j+2}}{4h^2}$ for $j = 0, 1, 2, ..., N^{-1}$
Again, we assume $u_{-1} = u_{N-1}$, $u_1 = u_{N+1}$, $u_{-2} = u_{N-2}$,..., etc.
Motivation: 0 $u(x_j + 2h) \approx u(x_j) + 2h u'(x_j) + 2h^2 u''(x_j)$
($0 + 0$): $u(x_j - 2h) \approx u(x_j) - 2h u'(x_j) + 2h^2 u''(x_j)$
This time, we approximate $\frac{du}{dx}$ as:
(**) $\frac{du}{dx}(x_j) \approx \frac{u_{j+1} - u_{j-1}}{2h}$

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Last time: an eigenvector of D and D. Claim: eikx is More specifically, $(De^{ikx})_j = \left(\frac{-\sin^2(kh)}{h^2}\right)e^{ikx}_j$ $\therefore De^{ikx}_j = \lambda_k^2 e^{ikx}_j$ $\left(\widetilde{D} e^{ihx}\right)_{j} = \left(\frac{i \sin lkh}{h}\right) e^{ikxj} \qquad \qquad \widetilde{D} e^{ikx} = \widetilde{\lambda}_{k} e^{ikx}$

Now,
$$a \frac{d^{2}u}{dx^{2}} + b \frac{du}{dx} = f$$
 can be discretized as:
 $a\widetilde{D}^{2}\widetilde{u} + b\widetilde{D}\widetilde{u} = \widetilde{f}$ subject to the periodic condition.
Recall: $\{ e^{i\frac{1}{k}x} \}_{k=0}^{N+1}$ is a basis for C^{N}
Again, let $\widetilde{u} = \sum_{k=0}^{N^{-1}} \widehat{u}_{k} e^{i\frac{1}{k}x}$ and $\widetilde{f} = \sum_{k=0}^{N^{-1}} \widehat{f}_{k} e^{i\frac{1}{k}x}$
Then: $a\widetilde{D}^{2}\left(\sum_{k=0}^{N^{-1}} \widehat{u}_{k} e^{i\frac{1}{k}x}\right) + b\widetilde{D}\left(\sum_{k=0}^{N^{-1}} \widehat{u}_{k} e^{i\frac{1}{k}x}\right) = \sum_{k=0}^{N^{-1}} \widehat{f}_{k} e^{i\frac{1}{k}x}$
 $\Longrightarrow \sum_{k=0}^{N^{-1}} \left(a\widetilde{\lambda}_{k}^{2} + b\widetilde{\lambda}_{k}\right) \widehat{u}_{k} e^{i\frac{1}{k}x} = \sum_{k=0}^{N^{-1}} \widehat{f}_{k} e^{i\frac{1}{k}x}$
Comparing coefficients, we get = $(a\widetilde{\lambda}_{k}^{2} + b\widetilde{\lambda}_{k}) \widehat{u}_{k} = \widehat{f}_{k}$ for $k=0,...,N^{-1}$
(algebraic equation)

$$\frac{Remark:}{a \tilde{\lambda}_{k}^{2} + b \tilde{\lambda}_{k}} = a \left(-\frac{\sin^{2} kh}{h^{2}} \right) + b \left(i \frac{\sin kh}{h} \right)$$

$$\therefore a \tilde{\lambda}_{k}^{2} + b \tilde{\lambda}_{k} = 0 \quad \text{if and only if } \sin kh = 0$$

$$\therefore k = \frac{2\pi}{N}$$

$$\therefore \sin kh = 0 \iff \sin \left(\frac{k 2\pi}{N} \right) = 0 \iff k = 0 \quad \text{or } k = \frac{N}{2}$$

$$(\text{for } k = 0, 1, 2, ..., N-1)$$

$$(assuming N is even)$$

For
$$k=0$$
 and $\frac{N}{2}$, $\tilde{\lambda}_{k}=0$. We set $\hat{u}_{0}=0=\hat{u}_{\frac{N}{2}}$.
In general, we set: $\hat{u}_{k}=\begin{cases} \hat{f}_{k}/(a\tilde{\lambda}_{k}^{2}+b\tilde{\lambda}_{k}) \text{ for } k\neq 0, \frac{N}{2}\\ 0 & \text{ for } k=0 \text{ or }\\ k=N_{2} \end{cases}$

 \vec{u} can be obtained by inverse DFT = $\begin{pmatrix} u_0 \\ \vdots \\ u_{N-1} \end{pmatrix} = A_w \begin{pmatrix} u_0 \\ \vdots \\ \vdots \\ u_{N-1} \end{pmatrix}$

<u>Answer</u>: How about the general solution? <u>Answer</u>: Examine N(aD²+bD), the null space.

Claim:
$$N(aD^{2} + bD) = span \{e^{i(0)X}, e^{i(M)X}\}$$

 $= span \{(i), (e^{i(M)X})\}$
Proof: $aD^{2} + bD$ has two eigenvectors whose eigenvalue is D.
These eigenvectors are $e^{i(0)X}$ and $e^{i(M)X}$.
 $N(aD^{2} + bD) = eigenspace of eigenvalue = 0$
 $= span \{e^{i(0)X}, e^{i(M)X}\}$.
Thus, general $sol = it^{X} = it + c_{1}e^{i(0)X} + c_{2}e^{i(M)X}$.
 $for some C_{1}$ and c_{2} .
 C_{1} and C_{2} can be determined by certain conditions (such as boundary conditions)

Main idea of numerical spectral method
Diff. egt
$$\rightarrow$$
 Algebraic egt \rightarrow solution \overline{u}
 $\begin{pmatrix} e.g. \\ d^{2u} = f \end{pmatrix}$ $\begin{pmatrix} \overline{Aw} \overline{f} \end{pmatrix}$ $Fast!$ $\begin{pmatrix} Aw \widehat{u} \end{pmatrix}$
 $Fast?$
Remark: To develop an efficient numerical spectral method,
 $we need to compute Aw \widehat{u}$ and $\overline{Aw} \overline{f}$ fast.
Computational cost for $Aw \widehat{u}$ is $O(n^2)$.
Goal: Reduce the computational cost to $O(n \log n)$
 $e.g. n = 2^{10}, n^2 = 2^{20}, n \log n = 10.2^{10} < 2^{14}$ $\therefore 2^6 = 64$ times faster.