3.2 Duality

3.2.1 Lagrangian and Dual Function

We consider a standard optimization problem (P):

$$\min f(x)$$
subject to $g_i(x) \le 0, i = 1, ..., h$

$$h_j(x) = 0, j = 1, ..., k$$

The optimal value p^* of (P) is called the primal optimal value.

Definition: (Lagrangian) The *Lagrangian* associated with the above problem is defined as

$$L(x, \lambda, \mu) = f(x) + \sum_{i=1}^{h} \lambda_i g_i(x) + \sum_{j=1}^{k} \mu_j h_j(x)$$

The vectors λ, μ are called the dual variables or Lagrange multipliers.

Definition: (Dual function) The dual function is defined as

$$q(\lambda, \mu) = \inf_{x} L(x, \lambda, \mu)$$

Note that q is always concave, being the pointwise infimum of affine functions.

Let p^* be the optimal value of (P). The dual function gives a lower bound on p^* .

Proposition: For all $\lambda \geq 0$ and μ , we have

$$q(\lambda, \mu) \le p^*$$

Proof. Let x be a feasible point. Then $g_i(x) \leq 0$ and $h_j(x) = 0$. Then

$$\sum_{i=1}^{h} \lambda_i g_i(x) + \sum_{j=1}^{k} \mu_j h_j(x) \le 0$$

Hence for all $\lambda \geq 0$ and μ ,

$$q(\lambda, \mu) \le L(x, \lambda, \mu) = f(x) + \sum_{i=1}^{h} \lambda_i g_i(x) + \sum_{j=1}^{h} \mu_j h_j(x) \le f(x)$$

Since this holds for all feasible points, we have $q(\lambda, \mu) \leq p^*$.

We next consider the dual problem.

Definition:(**Dual Problem**) The following optimization problem (D) is called the *dual problem* associated to (P):

$$\max q(\lambda, \mu)$$
 subject to $\lambda \ge 0$

A pair (λ, μ) such that $\lambda \geq 0$ and $q(\lambda, \mu) > -\infty$ is called *dual feasible*. A optimal solution (λ^*, μ^*) is called *dual optimal*.

Example: (Linear Program)

Consider a standard linear program (LP):

$$\min_{x \in \mathbb{R}^n} \langle c, x \rangle$$
 subject to $Ax = b$
$$x > 0$$

The Lagrangian is given by

$$L(x, \lambda, \mu) = c^{T} x - \sum_{i=1}^{n} \lambda_{i} x_{i} + \mu^{T} (Ax - b) = (A^{T} \mu + c - \lambda)^{T} x - b^{T} \mu$$

If $c+A^T\mu-\lambda\neq 0$, then $L(x,\lambda,\mu)$ is unbounded below. Hence the dual function is given by

$$q(\lambda, \mu) = \begin{cases} -b^T \mu & c + A^T \mu - \lambda = 0 \\ -\infty & \text{otherwise} \end{cases}$$

Therefore, the dual problem is given by

$$\max \ -b^T \mu$$
 subject to
$$A^T \mu + c - \lambda = 0$$

$$\lambda \ge 0$$

It can also be written in this form:

$$\max \ -b^T \mu$$
$$A^T + c\mu \ge 0$$

Example: Duality and Conjugate function

Consider the following optimization problem

$$\min f(x)$$
subject to $Ax \le b$

$$Cx = d$$

The dual function is

$$q(\lambda, \mu) = \inf_{x} (f(x) + \lambda^T (Ax - b) + \mu^T (Cx - d))$$
$$= -b^T \lambda - d^T \mu + \inf_{x} (f(x) + (A^T \lambda + C^T \mu)^T x)$$

Note that

$$\inf_{x} (f(x) + (A^T \lambda + C^T \mu)^T x) = -\sup_{x} (-(A^T \lambda + C^T \mu)^T x - f(x)) = -f^*(-(A^T \lambda + C^T \mu))$$

Hence, we have

$$q(\lambda, \mu) = -b^T \lambda - d^T \mu - f^*(-(A^T \lambda + C^T \mu))$$

3.2.2 Strong and Weak Duality

Let d^* be the optimal value of the dual problem. We have the following inequality.

Proposition:(Weak Duality) Let p^* be the primal optimal value and d^* be the dual optimal value. Then

$$d^* \le p^*$$

The difference $p^* - d^*$ is called the *duality gap*. If $p^* = d^*$, then we say that *strong duality* holds.

This leads us to ask the question when do strong duality holds. Such conditions are called constraint qualification. We will study one simple qualification: Slater's condition.

Consider a convex problem of the form:

min
$$f(x)$$

subject to $g_i(x) \le 0, i = 1,...,h$
 $Ax = b$

where f, g_i are convex.

Slater's Condition: There exists $x \in ri(D)$ such that

$$q_i(x) < 0, i = 1, ..., h, Ax = b$$

where $D = \operatorname{dom} f \cap (\cap_i \operatorname{dom} q_i)$.

Theorem:(Slater's Theorem) If the problem is convex and Slater's condition is satisfied, then strong duality holds.

3.2.3 Geometric Interpretation

Consider the following set

$$A := \{(u, v, t) | \exists x \ g_i(x) \le u_i, \ i = 1, ..., h, \ h_i(x) = v_i, \ j = 1, ..., k, \ f(x) \le t \}$$

We can show that A is convex if the problem is convex. Note that

$$p^* = \inf\{t | (0, 0, t) \in A\}$$

that is the lowest point where A intersects the 'vertical'-axis.

We can also interpret the dual function through this geometric setting:

$$q(\lambda, \mu) = \inf\{\langle (\lambda, \mu, 1), (u, v, t) \rangle | (u, v, t) \in A\}$$

For fixed (λ, μ) , we can define a hyperplane

$$\langle (\lambda, \mu, 1), (u, v, t) \rangle = q$$

Then $q(\lambda,\mu)$ is where a supporting hyperplane to A with 'slope' (λ,μ) intersects the 'vertical' axis.

Therefore, strong duality holds if and only if there is a nonvertical supporting hyperplane to A at $(0,0,p^*)$.

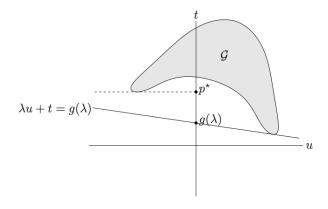


Figure 1: Geometric picture of the set G and dual function

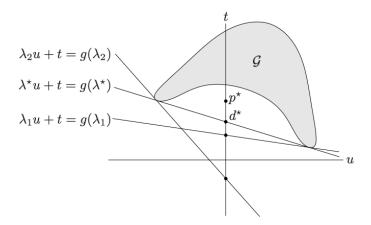


Figure 2: Primal and dual optimal value

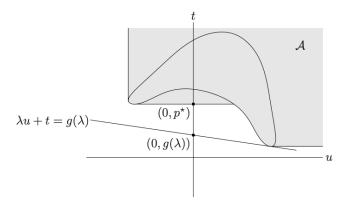


Figure 3: Geometric picture of the set A

Example: Consider the problem

$$\min_{x,y \ge 0} e^{-\sqrt{xy}}$$

subject to $x = 0$

The optimal value p^* is 1.

The dual function is given by

$$q(\lambda) = \inf_{x,y \ge 0} \{e^{-\sqrt{xy}} + \lambda x\} = \begin{cases} 0 & \lambda \ge 0 \\ -\infty & \lambda < 0 \end{cases}$$

Hence, the dual optimal value d^* is 0.

Therefore, the strong duality does not hold.

Note that Slater's Condition is not satisfied for this example.

3.3 KKT conditions

Let's consider the general convex problem again

$$\min f(x)$$
 subject to $g_i(x) \leq 0, \ i = 1, ..., h$
$$h_j(x) = 0, \ j = 1, ..., k$$

where are the functions are convex. We also assume that h_j are affine. Note that

$$\sup_{\lambda \ge 0, \mu} L(x, \lambda, \mu) = \begin{cases} f(x) & g_i(x) \le 0, \ h_j(x) = 0 \\ \infty & \text{otherwise} \end{cases}$$

Then $p^* = \inf_x \sup_{\lambda \geq 0, \mu} L(x, \lambda, \mu)$ On the other hand, $d^* = \sup_{\lambda \geq 0, \mu} \inf_x L(x, \lambda, \mu)$. Therefore, strong duality is equivalent to

$$\sup_{\lambda \geq 0, \mu} \inf_x L(x, \lambda, \mu) = \inf_x \sup_{\lambda \geq 0, \mu} L(x, \lambda, \mu)$$

Suppose strong duality holds. Let x^* be primal optimal and (λ^*, μ^*) be dual optimal. Then

$$\begin{split} f(x^*) &= q(\lambda^*, \mu^*) \\ &= \inf_x (f(x) + \sum_{i=1}^h \lambda_i^* g_i(x) + \sum_{j=1}^k \mu_j^* h_j(x)) \\ &\leq f(x^*) + \sum_{i=1}^h \lambda_i^* g_i(x^*) + \sum_{j=1}^k \mu_j^* h_j(x^*) \\ &\leq f(x^*) \end{split}$$

Therefore, we have equality for each line. In particular, we have

$$\sum_{i=1}^{h} \lambda_i^* g_i(x^*) = 0$$

Since each term is nonpositive, we have $\lambda_i^* g_i(x^*) = 0$ for all i.

This is called *complementary slackness*.

Suppose all the functions are also differentiable. Then since x^* minimize $L(x, \lambda^*, \mu^*)$, we have

$$\nabla_x L(x^*, \lambda^*, \mu^*) = 0$$

That is

$$\nabla f(x^*) + \sum_{i=1}^h \lambda_i^* \nabla g_i(x^*) + \sum_{j=1}^k \mu_j^* h_j(x^*) = 0$$

Combining with the complementary slackness condition, we have the following Karush-Kuhn-Tucker(KKT) condition:

$$\nabla f(x^*) + \sum_{i=1}^h \lambda_i^* \nabla g_i(x^*) + \sum_{j=1}^k \mu_j^* \nabla h_j(x^*) = 0$$

$$g_i(x^*) \le 0, \ i = 1, ..., h$$

$$h_j(x^*) = 0, \ j = 1, ..., k$$

$$\lambda_i^* \ge 0$$

$$\lambda_i^* g_i(x^*) = 0, \ i = 1, ..., h$$

Conversely, suppose $x^*, (\lambda^*, \mu^*)$ satisfy the KKT conditions. Since $L(x, \lambda^*, \mu^*)$ is convex in x and $\nabla_x L(x^*, \lambda^*, \mu^*) = 0$, then x^* minimizes $L(x, \lambda^*, \mu^*)$. Then

$$q(\lambda^*, \mu^*) = L(x^*, \lambda^*, \mu^*) = f(x^*) + \sum_{i=1}^h \lambda_i^* g_i(x^*) + \sum_{j=1}^h \mu_j^* h_j(x^*) = f(x^*)$$

Therefore, there is no duality gap and x^* , (λ^*, μ^*) are primal optimal and dual optimal respectively.

To conclude, we have the following optimal condition:

Theorem: Consider the convex problem (P). Suppose strong duality holds. Then x^* , (λ^*, μ^*) are primal and dual optimal if and only if x^* , (λ^*, μ^*) satisfy the KKT conditions.

Remark: If the functions are not differentiable, we can replace the first KKT condition by $0 \in \partial f(x^*) + \sum_{i=1}^h \lambda_i^* \partial g_i(x^*) + \sum_{j=1}^k \mu_j^* \partial h_j(x^*)$.

Example Consider the problem

$$\min x^2 + y^2$$
subject to $x + y = 1$

$$x, y \ge 0$$

The KKT condition can be written as

$$2x - \lambda_1 + \mu = 0$$
$$2y - \lambda_2 + \mu = 0$$
$$x + y = 1$$
$$x, y \ge 0$$
$$\lambda_1, \lambda_2 \ge 0$$
$$\lambda_1x = \lambda_2y = 0$$

By the first two conditions, we have,

$$\lambda_1 = 2x + \mu, \ \lambda_2 = 2y + \mu.$$

By the complementary slackness conditions, we have

$$2x^2 + \mu x = 0$$
, $2y^2 + \mu y = 0$.

Then x = 0 or $-\mu/2$, and y = 0 or $-\mu/2$.

We cannot have x, y both equal to 0 since otherwise $x + y = 0 \neq 1$.

Suppose exactly one of x, y is zero, say x.

Then $y = -\mu/2$. Since x + y = 1, then $\mu = -2$. But since x = 0, this implies that $\lambda_1 = \mu = -2 < 0$. This violates the dual feasibility.

Therefore, we must have x, y both nonzero. That is $x = y = -\mu/2$.

Since x + y = 1, we have $\mu = -1$. So x = y = 1/2 and $\lambda_1 = \lambda_2 = 0$ satisfy the KKT conditions. Therefore the global minimum is obtained at (1/2, 1/2).