

# Chapter S: Centrality in Social Network

## S.1 Centralities

One basic but essential measure in social network analysis is centrality. There are 4 common centrality measures which are degree centrality, betweenness centrality, closeness centrality, and eigenvector centrality. They have been employed to understand the roles of certain nodes in networks.

|                        |   |
|------------------------|---|
| Degree centrality      | Average degree of a vertex  |
| Betweenness centrality | Extent to which a particular vertex lies on the shortest path between other vertices  |
| Closeness centrality   | The average of the shortest distances to all other vertices   |
| Eigenvector centrality | A measure of the extent of which a vertex is connected to influential other vertices. Related concept is Google's Page Rank |

Let  $G = (V, E)$  be a simple connected  $(p, q)$ -graph. The degree centrality of a vertex is defined to be its degree. But for comparison, some article defines the *degree centrality* of a vertex  $x$  by

$$C_D(x) = \frac{\deg(x)}{p-1}.$$

This makes the scale between 0 and 1.

Degree centrality <https://www.youtube.com/watch?v=iiVeQkIELyc>

Suppose  $x, u, v \in V$  are distinct vertices in  $G$ . Let  $g_{u,v}$  denote the number of shortest path between  $u$  and  $v$ ; and  $g_{u,v}(x)$  denote the number of shortest path between  $u$  and  $v$  that pass through  $x$ .

The *betweenness centrality* of  $x$  is defined by

$$C_B(x) = \frac{1}{\binom{n-1}{2}} \sum_{\substack{u \neq v \\ x \notin \{u,v\}}} \frac{g_{u,v}(x)}{g_{u,v}}.$$

Betweenness centrality <https://www.youtube.com/watch?v=0CCrqr62TF7U&t=239s>

Betweenness centrality <https://www.youtube.com/watch?v=ptqt2zr9ZRE&t=1s>

Let  $d(u, v)$  denote the distance from  $u$  to  $v$ , where  $u, v \in V$ . For  $x \in V$ , the *closeness centrality* of  $x$  is defined by

$$C_C(x) = \frac{p-1}{\sum_{v \in V \setminus \{x\}} d(x, v)}.$$

Closeness centrality <https://www.youtube.com/watch?v=0unzqsPaPk8>

Eigenvector centrality: Score of a page is proportional to the sum of the scores of pages linked to it.

Let  $A$  be the adjacent matrix of a connected graph  $G$ . Since  $A$  is symmetric,  $A$  is diagonalizable over  $\mathbb{R}$ . From Perron-Frobenius Theorem, we know that  $A$  is nonnegative definite. Let  $\rho(A) = \rho$  be the largest eigenvalue which is called *Perron-Frobenius eigenvalue*. Then  $\rho$  is a simple eigenvalue, i.e., the eigenspace corresponding to  $\rho$  is of dimension one. Let  $\mathbf{v}$  be the eigenvector corresponding to  $\rho$ . Then all entries of  $\mathbf{v}$  are positive. So we may assume that  $\mathbf{v}$  is unit length. Note that,  $\mathbf{v}$  is the only unit eigenvector having this property.

The entry of  $\mathbf{v}$  corresponding to the vertex is called the *eigenvector centrality* of that vertex.

Beginner <https://www.youtube.com/watch?v=9vs1zSqd070>

Eigenvector centrality <https://www.youtube.com/watch?v=q8oBwS2wAQ>

## S.2 Page Rank

A concept related to eigenvector centrality is Google's Page Rank.

Following is the Page Rank Algorithm: Given a  $(p, q)$ -digraph  $\vec{G} = (V, E)$ , let

$$h_{u,v} = \begin{cases} \frac{1}{\deg^+(u)} & \text{if } (u, v) \in E; \\ 0 & \text{otherwise,} \end{cases}$$

for  $u, v \in V$ .

Let  $H = (h_{u,v})^T$  and  $\mathbf{v} = \frac{1}{p}\mathbf{1}$ , where  $\mathbf{1}$  is the  $p$ -vector whose entries are 1. Let  $\mathbf{v}_n = H^n \mathbf{v}$ . Then rank of the vertices follow the natural order of their corresponding entries of  $\mathbf{v}_n$  for some iteration  $n \geq 1$ , the largest entry corresponding to the highest rank.

Explanation <https://www.youtube.com/watch?v=MG0fIXfrT9A>

Page Rank Algorithm [https://www.youtube.com/watch?v=P8Kt6Abq\\_rM](https://www.youtube.com/watch?v=P8Kt6Abq_rM)

Original Formula <https://www.youtube.com/watch?v=pA1Q1myuScs>

Matrix Representation <https://www.youtube.com/watch?v=kSmQbVxq0Jc>