

Tutorial 1

1. Combinatorial games.

Recall that a game is called a combinatorial game if it satisfies the following axioms.

- (i) There are 2 players.
- (ii) There are finite many possible positions.
- (iii) The players take turns to make moves.
- (iv) Perfect information, i.e. both players know the rules and the possible moves of the other player.
- (v) The game ends in a finite number of moves without a draw.

A combinatorial game is said to be *impartial* if at any position, both players have the same possible moves.

Example 1. *Chess (if draws are not allowed) is a combinatorial game but it is not impartial.*

From now on, we only consider impartial combinatorial games.

Example 2. *The following games are (impartial) combinatorial games.*

- (i) *One-pile take away game.*
- (ii) *Two-pile take away game.*
- (iii) *Subtraction.*
- (iv) *Nim.*
- (v) *Nimber.*

Example 3. *The following games are not combinatorial games.*

- (i) *Poker.*

(ii) *Rock-paper-scissors.*

(iii) *Tic-tac-toe.*

Note. *In poker, there is no perfect information. The second game does not satisfy axiom (iii). In a game of tic-tac-toe, it is possible to get a draw.*

2. Solving combinatorial games.

Winning strategy: a strategy of a player that guarantees a win.

The following result is our starting point of studying combinatorial games.

Theorem 1. *(Zermelo). In any combinatorial game, exactly one of the players has a winning strategy.*

Now we may understand the problem of *solving a combinatorial game* as determining which player has a winning strategy at a given position. To do this, we need to introduce two core concepts.

N-positions and P-positions.

N-position: a position at which the player who is about to move has a winning strategy.

P-position: a position at which the player who just moved has a winning strategy.

Characterization of N-positions and P-positions:

- (i) All terminal positions are P-positions.
- (ii) From every P-position, any move can only reach an N-position.
- (iii) From every N-position, there exists at least one move to a P-position.

Note. *1. By the above characterization, the player who reaches a P-position has a winning strategy, i.e. keeping staying in P-positions. We call a move from an N-position to a P-position a winning move.*

2. The above characterization also provides an algorithm to find all P-positions, i.e. backwards induction.

Exercise 1. (Subtraction game). There is a pile of chips on the table. In each turn, a player removes 2 or 3 chips. The game ends when there are no possible moves and the player who makes the last move wins.

(i) Find all P-positions.

(ii) Find all winning moves from the position that there are 2019 chips.

Solution: (i). The set of terminal position is $\{0, 1\}$. By the characterization of N-positions and P-positions, we have by backwards induction,

0	1	2	3	4	5	6	7	8	9	10	11...
P	P	N	N	N	P	P	N	N	N	P	P...

Let

$$\mathcal{P} = \{k \in \mathbb{N} : k \equiv 0 \text{ or } 1 \pmod{5}\}.$$

We claim that \mathcal{P} is the set of all P-positions. Proof of the claim: (1). Clearly the terminal positions 0 and 1 are in \mathcal{P} . (2). If $k \in \mathcal{P}$, then we have $k - 2 \equiv 3$ or $4 \pmod{5}$ and $k - 3 \equiv 2$ or $3 \pmod{5}$, hence $k - 2 \notin \mathcal{P}$ and $k - 3 \notin \mathcal{P}$. (3). If $k \notin \mathcal{P}$, then $k \equiv 2 \pmod{5}$ or $k \equiv 3 \pmod{5}$ or $k \equiv 4 \pmod{5}$, in each case, we have either $k - 2 \equiv 0$ or $1 \pmod{5}$ or $k - 3 \equiv 0$ or $1 \pmod{5}$, hence either $k - 2 \in \mathcal{P}$ or $k - 3 \in \mathcal{P}$. By the characterization of P-positions, we complete the proof.

(ii). Clearly $2019 \equiv 4 \pmod{5}$ is an N-position and the only winning move from 2019 is removing 3 chips to reach 2016.

3. Nim.

Model: suppose there are n piles of chips on the table. A move consists of choosing one pile and removing arbitrary positive number of chips from this pile. The player who makes the last move wins.

Let \mathcal{P} denote the set of P-positions.

Theorem 2. $\mathcal{P} = \{(x_1, \dots, x_n) : x_1 \oplus \dots \oplus x_n = 0\}$.

Exercise 2. Consider 3-pile nim. Find all winning moves from position $(29, 20, 15)$.

Solution: Note that

$$29 \oplus 20 \oplus 15 = \begin{array}{r} (1, 1, 1, 0, 1)_2 \\ (1, 0, 1, 0, 0)_2 \\ (0, 1, 1, 1, 1)_2 \\ \hline (0, 0, 1, 1, 0)_2 = 6 \end{array} .$$

Hence $(29, 20, 15)$ is an N-position. All winning moves are $(1, 1, 1, 0, 1)_2 \rightarrow (1, 1, 0, 1, 1)_2 = 27$, $(1, 0, 1, 0, 0) \rightarrow (1, 0, 0, 1, 0)_2 = 18$, $(0, 1, 1, 1, 1)_2 \rightarrow (0, 1, 0, 0, 1)_2 = 9$. In other words, the winning moves are removing 2 chips from the first pile, or removing 2 chips from the second pile, or removing 6 chips from the third pile.